

Numerical Linear Algebra
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Lecture - 27
Stability of Numerical Algorithms- I

Hello friends. I welcome you to my lecture on Stability of Numerical Algorithms. There will be 2 lectures on this topic, this is the first lecture on stability of numerical algorithms.

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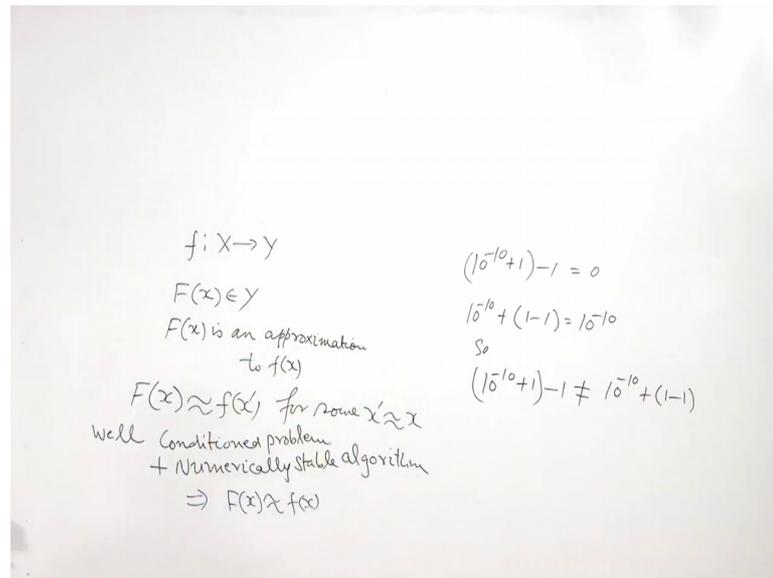
We know that a problem is a map $f: X \rightarrow Y$ from a normed vector space of data to a normed vector space Y of solutions. A numerical algorithm is a procedure which calculates $F(x) \in Y$, an approximation to $f(x)$. A numerical algorithm does not necessarily have to be finite. Some algorithms converge to the true solution in the limit.

A numerical algorithm is called numerically stable if it is not sensitive to small round of errors because round of errors are inevitably introduced at some stage in the numerical computations. Thus, an algorithm is stable if $F(x) \approx f(x')$ for some $x' \approx x$, which means that a stable algorithms computes “nearly the right answer” to “nearly the right question”.

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We have seen that a problem is a function f from X into Y , where X and Y are normed vector spaces X is the normed space normed vector space of data and Y is the normed vector space of solutions. A numerical algorithm is a procedure which calculates $F x$, capital $F x$ belonging to by f is the function from X into Y , where X and Y are non vector spaces and numerical algorithm is a procedure which calculates capital $F x$ belonging to y , capital $F x$ belonging to y which is an approximation to the actual value $F x$.

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So, this $F(x)$ is then approximation to $f(x)$. A numerical algorithm does not necessarily have to be finite. Some algorithms converge to the 2 solution in the limiting case. A numerical algorithm is called numerically stable if it is not sensitive to small round of errors because round of errors are inevitably introduced at some stage in the numerical computations. Thus an algorithm is said to be stable if $f(x)$ is approximately equal to $F(x)$ is approximately equal to, $F(x')$ for some $x' \approx x$, which means that a stable algorithm computes nearly the right answer to nearly the right question.

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Let us note the contrast between conditioning and stability:

- conditioning applies to problems
- stability applies to algorithms.

Thus, stability + good conditioning \Rightarrow accuracy. If a stable algorithm is applied to a well conditioned problem then $F(x) \approx f(x)$.

Conversely, if a problem is ill-conditioned then an accurate solution may not be possible or meaningful.

Floating point arithmetic is stable for computing sums, products, quotients and differences of two numbers.

Floating point operations are not associative.

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Let us note the contrast between the conditioning and stability. Conditioning as we have seen applies to problems and stability applies to algorithms. Thus we can say that stability plus good conditioning implies the accuracy.

If a stable algorithm is applied to a well conditioned problem then capital F x is approximately equal to. So, conditioning well conditioned problem if we have well conditioned problem if we have a well conditional, well conditional problem, then well condition problem plus numerically stable algorithm implies that F x is approximately equal to f x. Now, conversely if a problem is ill conditioned then an accurate solution may not be possible are meaningful.

Floating point arithmetic is stable for computing sums, products, quotients and differences of 2 numbers. Floating point operations however are not associative as we can see from this example.

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Example: Let us assume 10 digit precision then

$$(10^{-10} + 1) - 1 = 0 \text{ but } 10^{-10} + (1 - 1) = 10^{-10}.$$

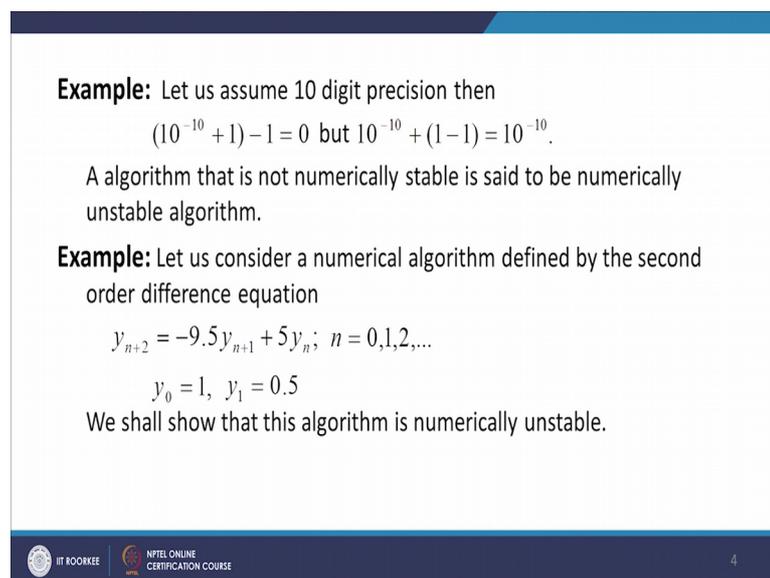
A algorithm that is not numerically stable is said to be numerically unstable algorithm.

Example: Let us consider a numerical algorithm defined by the second order difference equation

$$y_{n+2} = -9.5y_{n+1} + 5y_n; \quad n = 0, 1, 2, \dots$$

$$y_0 = 1, \quad y_1 = 0.5$$

We shall show that this algorithm is numerically unstable.



Let us assume 10 digit precision then 10 to the power minus 10 plus 1 minus 1 is equal to 0 in the 10 digit precision. But 10 to the power minus 10 plus 1 minus one is equal to 10 to the power minus 10 so that means, that the floating point arithmetic is not associative, 10 to the power minus 10 plus 1 minus 1 this equal to 0 in 1 digit consider 10 digit precision, but 10 to the power minus 10 plus 1 minus 1 is equal to 10 to the power minus 10. So, this is not equal to and therefore, the floating point arithmetic is not associative.

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Example: Let us assume 10 digit precision then
 $(10^{-10} + 1) - 1 = 0$ but $10^{-10} + (1 - 1) = 10^{-10}$.

A algorithm that is not numerically stable is said to be numerically unstable algorithm.

Example: Let us consider a numerical algorithm defined by the second order difference equation

$$y_{n+2} = -9.5y_{n+1} + 5y_n; n = 0, 1, 2, \dots$$
$$y_0 = 1, y_1 = 0.5$$

We shall show that this algorithm is numerically unstable.

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An algorithm that is not numerically stable is said to be numerically unstable algorithm. Let us consider a numerical algorithm defined by the second order different equation. The difference equation we have considered is y_{n+2} equal to minus 9.5 y_{n+1} plus 5 y_n where n value 0, 1, 2, 3 and so on and we are given that y_0 is equal to 1 and y_1 is equal to 0.5 we shall see that this algorithm is numerically unstable. So, y_{n+2} is equal to minus 9.5 y_{n+1} plus 5 y_n , n is equal to 0, 1, 2 and so on; y_0 is equal to 1, y_1 equal to 0.5.

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Thus, we get $y_n = \left(\frac{1}{2}\right)^n, n = 0, 1, 2, \dots$

$$y_{n+2} = -9.5y_{n+1} + 5y_n, n = 0, 1, 2, \dots$$
$$y_0 = 1, y_1 = 0.5$$

we may write

$$y_{n+2} + 9.5y_{n+1} - 5y_n = 0$$
$$\propto p(E)y_n = 0$$

where $p(E) = E^2 + 9.5E - 5$
 $E(y(x)) = y(x+h)$

$$y_n = A(0.5)^n + B(-10)^n$$

Putting $n=0$, we get

$$y_0 = A + B$$
$$\Rightarrow A + B = 1 \Rightarrow A = (1 - B)$$

Putting $n=1$, we get

$$y_1 = A(0.5) + B(-10)$$
$$0.5 = A(0.5) + B(-10)$$
$$\Rightarrow 0.5(-A) = -10B$$
$$\frac{1}{2}B = -10B \Rightarrow B = 0$$

Hence $A = 1$

Now, let us see how it is numerically unstable let us consider a machine where beta is equal to 10s so that means, the decimal system we are working and 3 digit rounding arithmetic; that means, t is equal to 3. The given equation can be rewritten as $y_{n+2} + 9.5y_{n+1} - 5y_n = 0$ or equivalently we can write it as $p(E)y_n = 0$, $p(E)$, where $p(E) = E^2 + 9.5E - 5$ and y_{E^k} mean where y_{E^k} mean it is the shifting operator. So, $E y_x$, $E y_x$ is equal to y_{x+h} . So, E is the shifting operator here.

Now, when we want to solve a difference equation we write the corresponding auxiliary equation. So, here p is equal to $E^2 + 9.5E - 5$. So, we have the auxiliary equation as $p(\lambda) = 0$, that is $2\lambda^2 + 19\lambda - 10 = 0$ when you factorize this equation, the factors are $2\lambda - 1$ into $\lambda + 10$ equal to 0 which gives us the, which give us the values of λ as 0.5 and -10 .

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The auxiliary equation $p(\lambda) = 0$

$$\Rightarrow 2\lambda^2 + 19\lambda - 10 = 0$$

$$(2\lambda - 1)(\lambda + 10) = 0$$

$$\Rightarrow \lambda = 0.5, -10.$$

Hence, the general solution of the given second order homogeneous difference equation is

$$y_n = A(0.5)^n + B(-10)^n. \quad \dots\dots\dots(1)$$

Since $y_0 = 1$, we have $A + B = 1$

$$y_1 = 0.5 \Rightarrow \frac{1}{2} = A\left(\frac{1}{2}\right) + B(-10).$$

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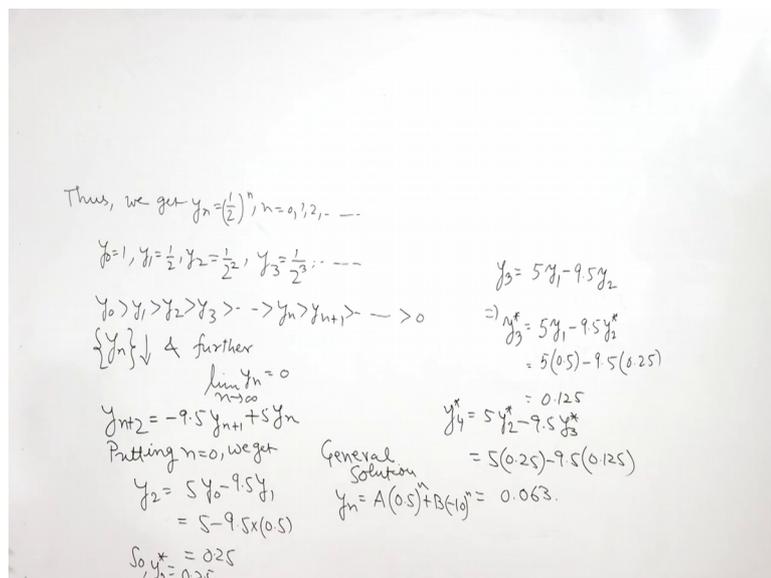
Hence the general solution of the given second order homogeneous difference equation is equal to $y_n = A(0.5)^n + B(-10)^n$. So, we have now in order to calculate the constants A and B we will make use of the given conditions $y_0 = 1$, $y_1 = 0.5$. So, putting $n = 0$ we get $y_0 = A + B$ and $y_0 = 1$, which gives you $A + B = 1$.

Now, let us put $n = 1$. So, putting $n = 1$ we get $y_1 = A(0.5) + B(-10)$, y_1 is given to be 0.5 , so $0.5 = A(0.5) + B(-10)$.

Now, $A + B = A$ gives you $B = 1 - A$ or you can say $A = 1 - B$, $A = 1 - B$. Now, what do we notice here? We can write it as $0.5 \times (1 - A) = -10B$, $1 - A = -20B$, $1 - A = 20B$. So, we write $0.5B = -10B$ which implies that $B = 0$ and when $B = 0$, $A = 1$, hence $A = 1$. And thus we get $y_n = \left(\frac{1}{2}\right)^n$ which is the general solution of the given second order homogeneous difference equation.

Now, from here we can note the following. From the sequence y_n we can see that as n increases y_n decreases. $y_0 = 1$, $y_1 = \frac{1}{2}$, $y_2 = \frac{1}{4}$, $y_3 = \frac{1}{8}$ and so on.

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So, $y_0 > y_1 > y_2 > \dots$ and $y_n > y_{n+1} > \dots$ and so on. All these values are greater than 0.

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$\{y_n\}$ is a strictly monotonically decreasing sequence of positive terms such that

$$\lim_{n \rightarrow \infty} y_n = 0.$$

Now, let us calculate the approximate solution using the difference equation and the machine arithmetic we get

$$y_2^* = 0.25, y_3^* = 0.125, y_4^* = 0.063$$
$$y_5^* = 0.027, y_6^* = 0.059, y_7^* = -0.426$$
$$y_8^* = 4.34, y_9^* = -43.4, y_{10}^* = 434.0\dots$$

The obtained approximate iterates are very absurd.

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Now, so, y_n is a strictly monotonically decreasing sequence of positive terms and also we see that, y_n is strictly monotonically decreasing sequence and further limit n tends to infinity y_n is equal to 0, 1 by 2 raise to the power n as n goes to infinity is equal to 0. Now, let us calculate the approximate solution using the difference equation and the machine arithmetic where we have assume that β equal to 10 and t equal to 3, decimal system having 3 digit rounding arithmetic.

So, our scheme is y_{n+2} is equal to $-9.5 y_{n+1} + 5 y_n$. So, these our scheme y_{n+2} equal to $5 y_n - 9.5 y_{n+1}$; y_0 and y_1 we are given, so we can start from y_2 . If you put n equal to 0 in this scheme putting n equal to 0 what do we get? y_2 equal to $5 y_0 - 9.5 y_1$. Now, y_0 is given to be equal to 1. So, we get 5 minus 9.5 into y_1 is given to be 0.5.

So, when you calculate this it comes out to be 0.25, we denote it by y_2^* . So, this, so we have computed values the values which are computed from the given difference equation we denoted by y_n^* . So, y_2^* is 0.25. Again we put n equal to 1 here we get y_3 equal to -9.5 in to y_2 plus $5 y_1$. So, y_3 which is equal to $5 y_1 - 9.5 y_2$ we are writing the computed values as y_n^* . So, this gives you y_3^* is equal to $5 y_1 - 9.5 y_2^*$ and $5 y_1$, $5 y_1$ is equal to 5 into 0.5 minus 9.5 into 0.25. So, when you compute this y_3^* comes out to be 0.125.

In this way we continue find y_4^* , y_4^* is equal to $5 y_2$; when we put n equal to 2 here. So, $y_2 y_2$ will become y_2^* minus $9.5 y_3^*$. So, putting the value of y_2^* as 0.25 we get 5 into 0.25 minus 9.5 into 0.125 which comes out to be 0.063 . So, like this we continue. And then we arrive at y_5^* , y_5^* is 0.027 y_6^* is 0.059 , y_7^* is minus 0.426 .

You can see now there is fluctuation of sine, earlier it has positive now it has become negative then we get y_8^* as 4 minus 34 ; 4.34 y_9^* is minus 43.4 . So, the values of y_n^* have started fluctuating starting with y_6^* and then y_{10}^* we can find y_{10}^* is very big 434.0 . So, the obtained approximate iterates are very absurd from the values of y_n^* we notice that they are absurd there is the values are absurd.

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While the actual iterates $\{y_n\}$ are monotonically decreasing and $\lim_{n \rightarrow \infty} y_n = 0$, the iterates $\{y_n^*\}$, as n gets large, oscillates and $|y_n^*| \rightarrow \infty$, as $n \rightarrow \infty$. The reason is that in the general solution (1) of the difference equation, the root $\lambda = -10$ also has a significant role to play. The round off error starts from the iterate $\{y_4^*\}$ onwards and is amplified by the multiplication factor -9.5 in each subsequent iteration. As a result, the round off errors grow uncontrollably and the computed iterates diverge from the actual iterates. A complete loss of significance of digits occurs for $\{y_8^*\}$ and it gets worse for higher values of n .

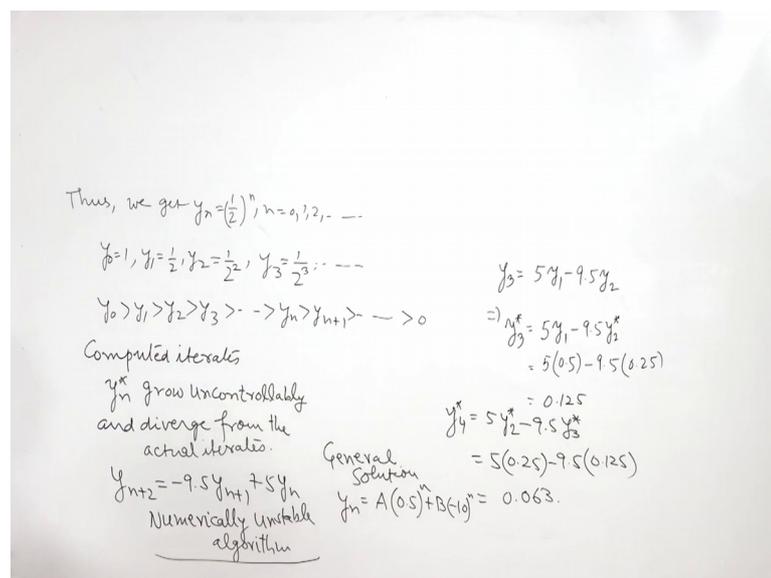
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The actual iterates y_n we have seen are monotonically decreasing, y_n is equal to 1 by 2 to the power n they are monotonically decreasing and y_n tends to 0 as n goes to infinity, while the iterates y_n^* as n gets large oscillates as we have seen and modulus of y_n^* tends to infinity. We can see that because the value of y_{10}^* has become 434 when you calculate y_{11}^* it will be much larger. So, n be the negative sign because they are fluctuating, but then the modulus of y_n^* , so numerically the y_n^* is tending to infinity when we take the modulus of y_n^* . So, it tends to infinity as n goes to infinity.

The reason is that in the general solution we have seen the general solution of the different equations general solution was y_n equal to A into 0.5 to the power n plus B into 10 to the power n . So, this the root λ equal to 10 has a significant role to play, but when what we do here. So, because of this 10 to the power n when we calculate y_n it comes out to be 1 by 2 to the power n then we get a monotonically decreasing sequence which can decrease to 0 . So, this 10 has a significant role. The round of error starts from the iterate y_4 onwards you can see y_4 is 0.063 and then y_5 is 0.027 .

So, y_2 that is y_2 star y_3 star 0.125 then y_4 star is 0.063 , here itself because of β equal to 3 the round of error starts taking place. And so, it gets, it in the next iterate y_5 star and so on it is amplified, because of the multiplication factor 9.5 . When you calculate y_5 star you will multiply by y_4 star by 9.5 . So, it gets amplified because of this factor in each subsequent iteration as a result the round off errors grow uncontrollably and the computed iterates diverge from the actual iterates. So, what will happen is that. So, as a result of this multiplication of 9.5 the computed iterates y_n star they grow uncontrollably, they grow uncontrollably and the computed iterates and diverge and diverge from the actual iterates.

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Now, we can notice one more thing, a complete loss of significance of digits occurs for y_8 star let us go to y_8 star, y_8 star you can see is 4.34 . Earlier if you look at the values of

the iterates they are 0.25, then 0.125, then 0.063, then 0.027, then 0.059, then minus 0.426. So, we had 0 point some value. Now here it y_8 becomes 4.34. So, there is a complete loss of significance of digits for y_8 and it gets worse for higher values of n because then we are multiplying them by minus 9.5 that that factor and so the error gets amplified the round of error gets amplified and the computed iterates get worse for higher values of n .

So, this is we can say this scheme are this numerical algorithm is n unstable algorithm. So, we have y_{n+1} , y_{n+2} equal to minus 9.5 y_{n+1} plus 5 times y_n it is unstable, numerically unstable algorithm. With this I would like to conclude my lecture.

Thank you very much for your attention.