

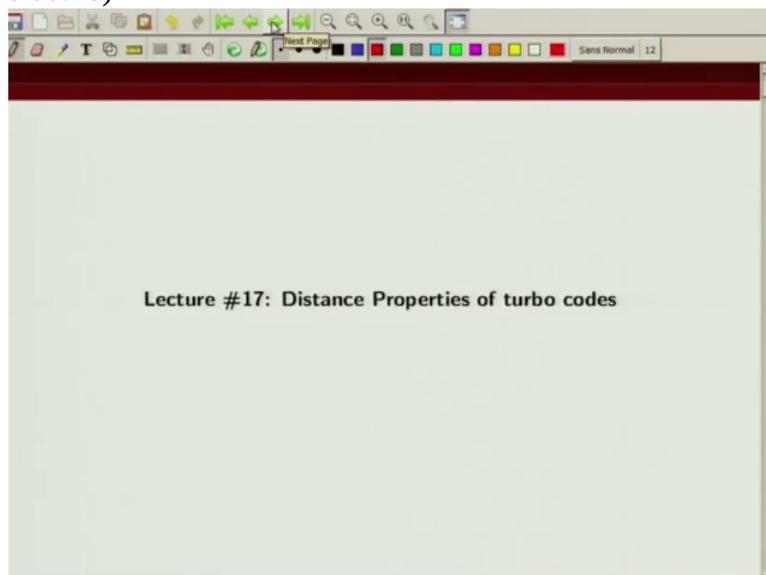
An Introduction to Coding Theory
Professor Adrish Banerji
Department of Electrical Engineering
Indian Institute of Technology, Kanpur
Module 08
Lecture Number 39
Distance Properties of Turbo Codes

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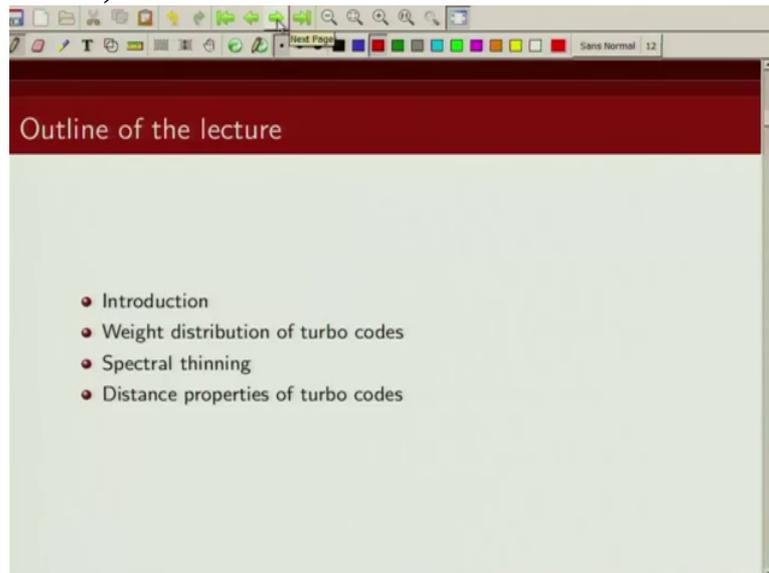
Welcome to the course on An Introduction to Coding Theory. I am Adrish Banerji from I I T Kanpur.

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In today's lecture we are going to talk about distance properties of turbo codes.

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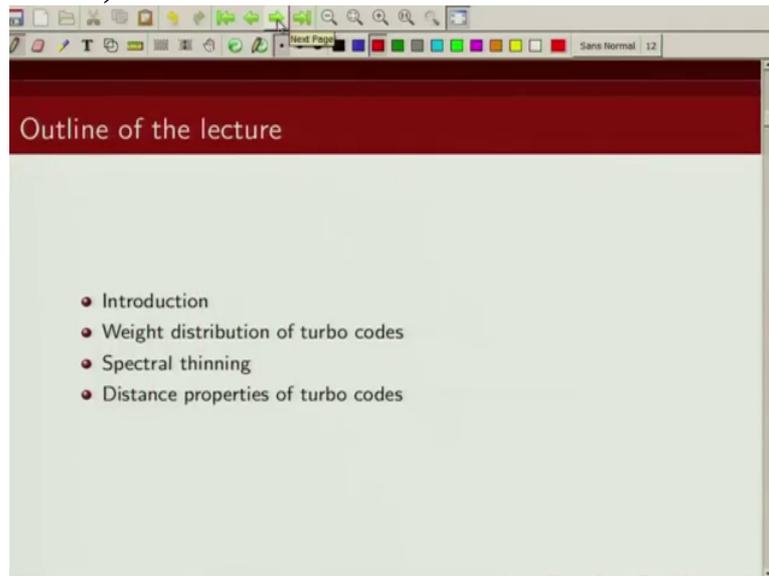
So this is the outline of today's talk. We will first give an example of weight distribution of a convolutional code

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and we will compare it with the weight distribution of a concatenated code. And we will see what effect interleaver has on the weight distribution of a concatenated code. Then we will introduce this concept of

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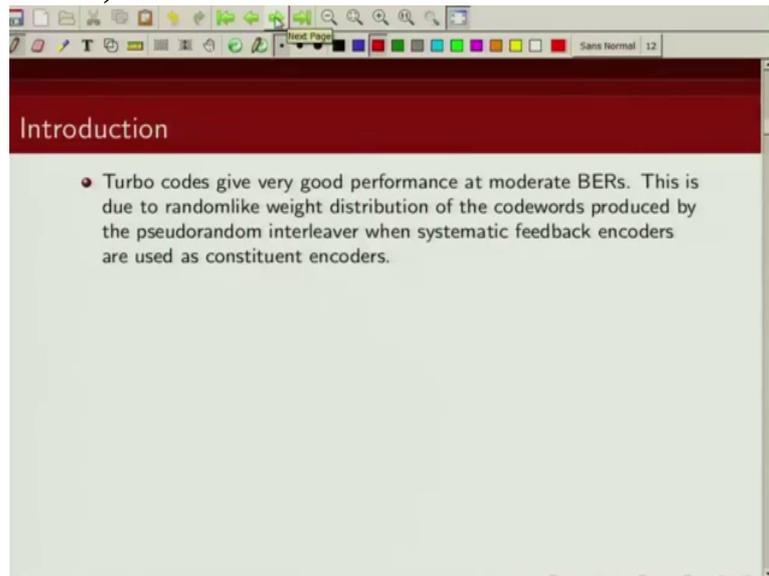
spectral thing which is visible in turbo code if we use feedback encoders and finally we will show how we can make use of this concept of uniform interleaver to derive

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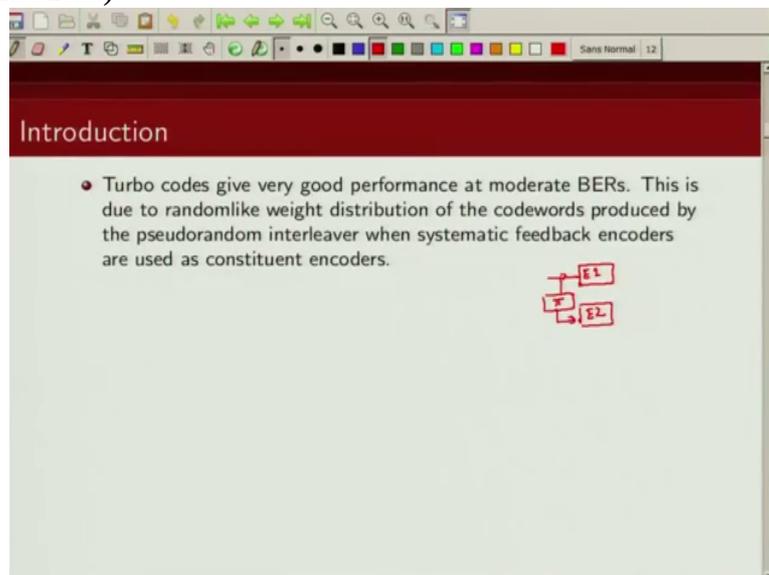
expression for average weight enumerating function of turbo code.

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So as we know turbo codes have very good performance at low signal to noise ratio and at moderate b e r and this is due to its random like weight distribution which is produced when we consider, pseudorandom interleavers combined with systematic feedback encoders. If you recall a turbo code is a parallel concatenation of two consequent encoders. So this is encoder 1, encoder 2. Information bit enters encoder 1 and the interleaved version of information enters

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encoder 2. These are the parity sequence coming out of first encoder and second encoder and then this is systematic, so typically you have this example of rate one third parallel concatenated code.

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The screenshot shows a presentation slide with a red header titled "Introduction". The main content area is white and contains a bullet point: "• Turbo codes give very good performance at moderate BERs. This is due to randomlike weight distribution of the codewords produced by the pseudorandom interleaver when systematic feedback encoders are used as constituent encoders." To the right of the text, there is a handwritten diagram in red ink. It shows two parallel encoder blocks, labeled E1 and E2, each with a feedback loop. The input to the top block is labeled 'x' and the output is labeled 'y'. The input to the bottom block is labeled 'x' and the output is labeled 'y'. To the right of the diagram, the text $R = \frac{1}{3} PCC$ is written.

Now we will try

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The screenshot shows a presentation slide with a red header titled "Introduction". The main content area is white and contains two bullet points: "• Turbo codes give very good performance at moderate BERs. This is due to randomlike weight distribution of the codewords produced by the pseudorandom interleaver when systematic feedback encoders are used as constituent encoders." and "• Let us consider some examples to understand this feature."

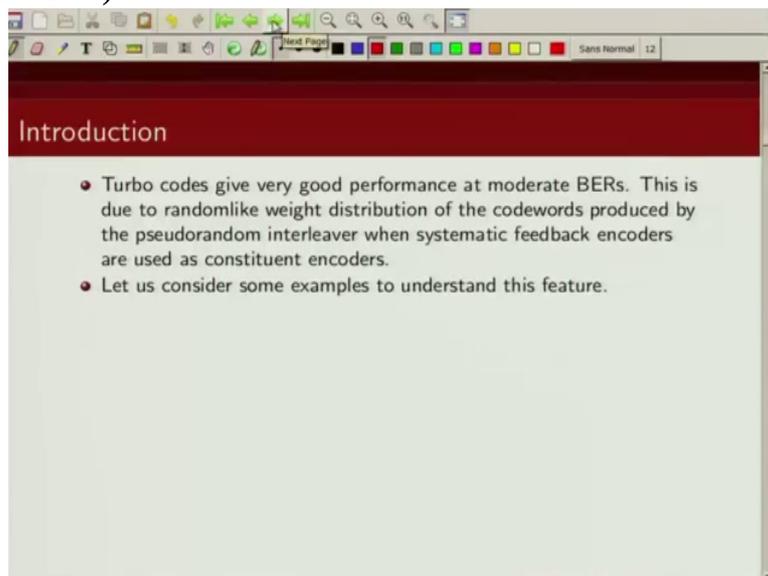
to see, as I said,

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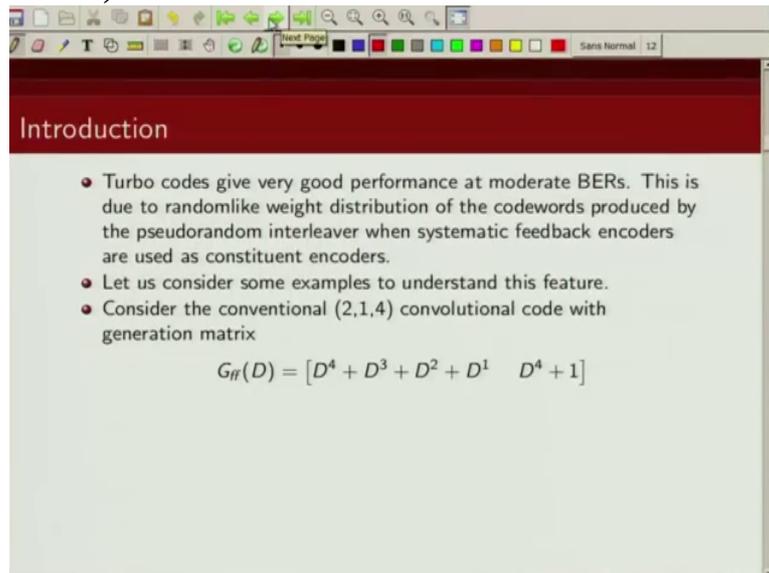
we will try to look at the distance spectrum of the concatenated code and see how it is different from a

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convolutional code.

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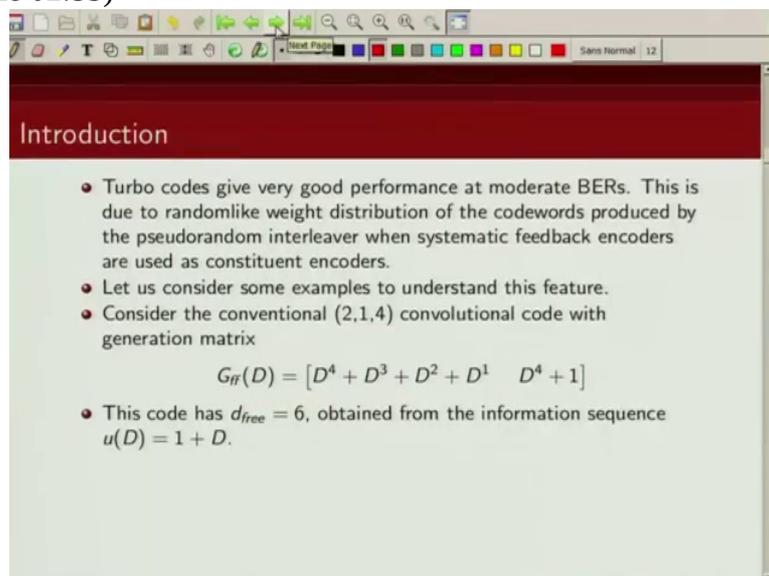
The slide is titled "Introduction" and contains the following text:

- Turbo codes give very good performance at moderate BERs. This is due to randomlike weight distribution of the codewords produced by the pseudorandom interleaver when systematic feedback encoders are used as constituent encoders.
- Let us consider some examples to understand this feature.
- Consider the conventional (2,1,4) convolutional code with generation matrix

$$G_{ff}(D) = [D^4 + D^3 + D^2 + D^1 \quad D^4 + 1]$$

So let us start with an example. So we are considering a rate one half, rate half convolutional code with memory 4 whose generator matrix is given by this. So this is a feed forward encoder. Now

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$$G_{ff}(D) = [D^4 + D^3 + D^2 + D^1 \quad D^4 + 1]$$

- This code has $d_{free} = 6$, obtained from the information sequence $u(D) = 1 + D$.

this code has a free distance of 6 which is obtained when we feed information of the form u equal to 1 plus d . So this is of the form, input sequence is 11 and all zeroes. Now we already know how to find

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Introduction

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- This code has $d_{free} = 6$, obtained from the information sequence $u(D) = 1 + D$. 11000...

a d_{free} of a convolutional code. Given a generator matrix you first have to find out its state diagram and then you draw modified state diagram and from there you can calculate weight enumerating function, input output weight enumerating function, so conditional weight enumerating function, so all those things we have already discussed in our earlier lecture. So you know how, given a generator matrix how to find out weight distribution of a convolutional code.

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Introduction

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- Let us consider some examples to understand this feature.
- Consider the conventional (2,1,4) convolutional code with generation matrix

$$G_{ff}(D) = [D^4 + D^3 + D^2 + D^1 \quad D^4 + 1]$$
- This code has $d_{free} = 6$, obtained from the information sequence $u(D) = 1 + D$.
- Now consider the equivalent (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_{fb}(D) = \left[1 \quad \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \right]$$

Now let us consider an equivalent convolutional encoder

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Introduction

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whose convolutional matrix is given by this. Now note this is equivalent to this feed forward

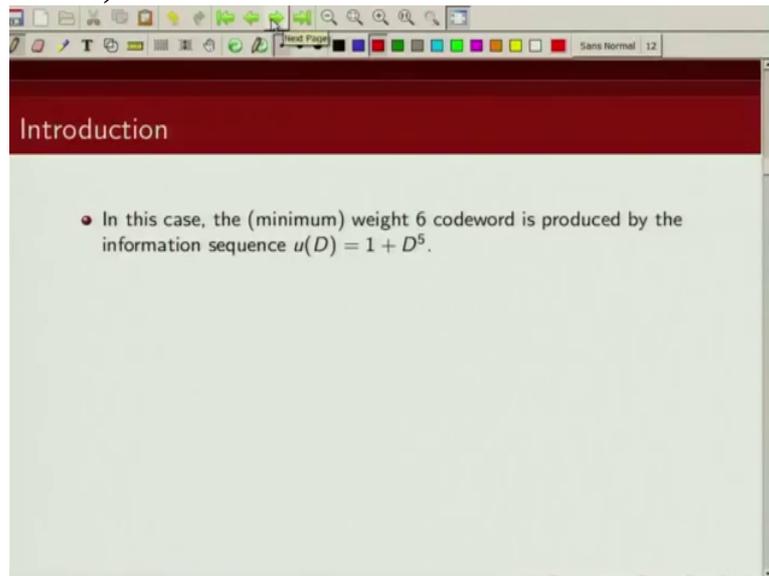
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Introduction

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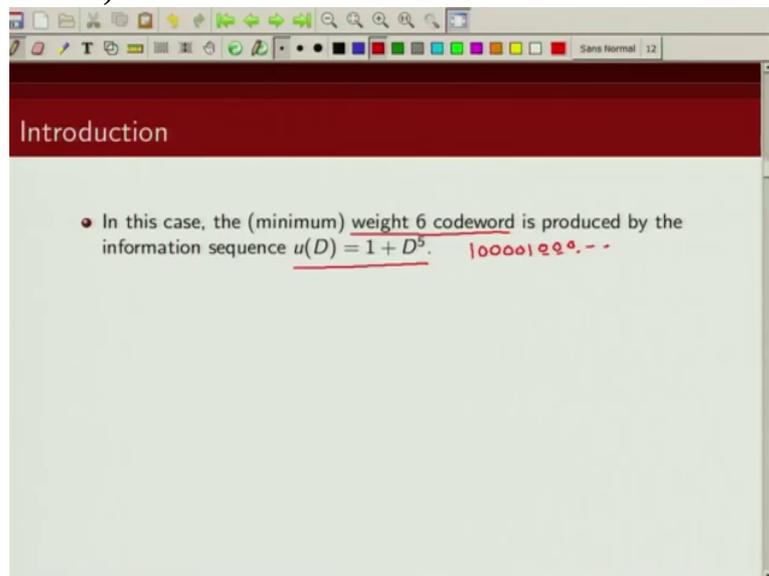
generator matrix, so this is a feedback encoder, and this is a systematic encoder. Now

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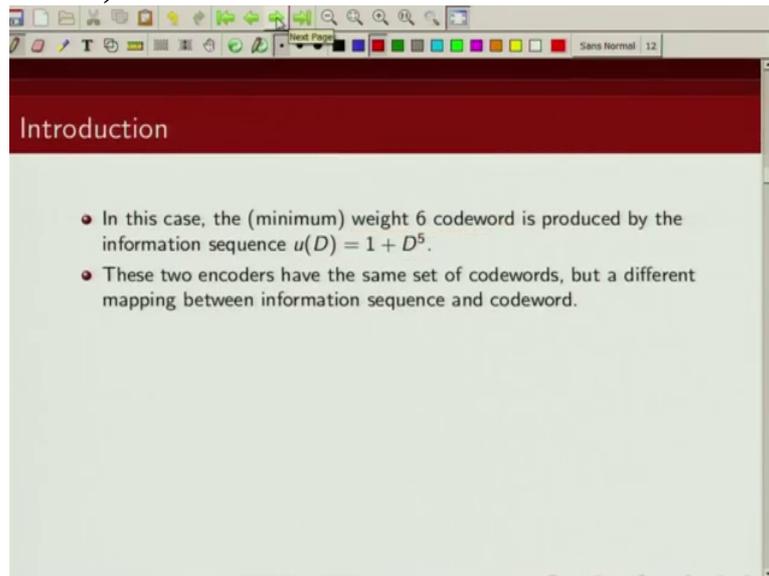


in this case of a systematic feedback encoder, the minimum weight codeword is again 6 and which is produced by input of the form 1 plus d^5 , so this is 1 0 0 0 0 1, so this is 1 d^2 square, d^3 cube, d^4 four d^5 five. So this is produced by sequence of this form.

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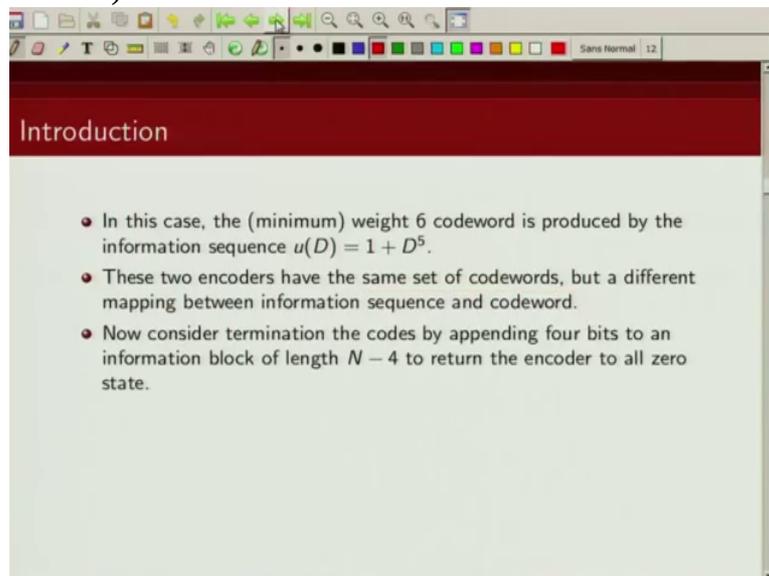


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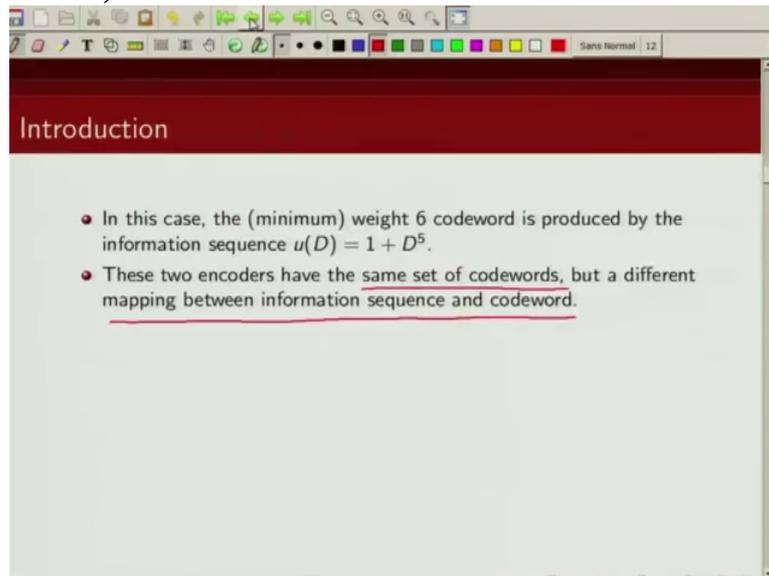
Now since these two generator matrix are equivalent, they will produce the same set of codewords. The only difference is the mapping between the information sequence and codeword will be different.

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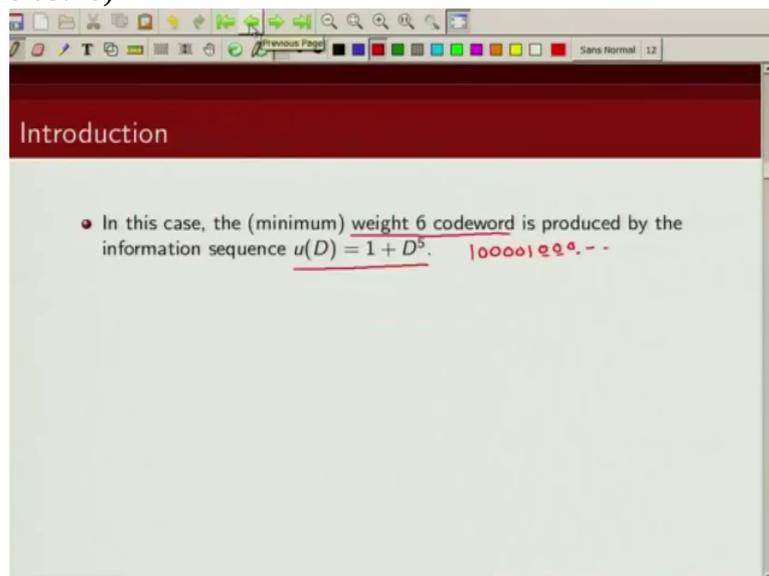
Now let us consider a terminated convolutional code. So let's say our overall information blocks we want is n . So in this example, since the memory

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here is 4,

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you can see

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- Consider the conventional (2,1,4) convolutional code with generation matrix
$$G_{ff}(D) = \begin{bmatrix} D^4 + D^3 + D^2 + D^1 & D^4 + 1 \end{bmatrix}$$
- This code has $d_{free} = 6$, obtained from the information sequence $u(D) = 1 + D$.
- Now consider the equivalent (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix
$$G_{fb}(D) = \begin{bmatrix} 1 & \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \end{bmatrix}$$

this is memory is 4, Ok. This is rate one half memory 4 code. So if we want to terminate this encoder we would require 4

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The slide is titled "Introduction" and contains the following text:

- In this case, the (minimum) weight 6 codeword is produced by the information sequence $u(D) = 1 + D^5$.
- These two encoders have the same set of codewords, but a different mapping between information sequence and codeword.
- Now consider termination the codes by appending four bits to an information block of length $N - 4$ to return the encoder to all zero state.

termination bits. So let us consider information block of length n minus 4 and to this we will add 4 termination bits. So overall codeword will be, because it is a rate half code, overall length of the codeword would be 2 times n . So this is information sequence n minus 4 and plus we add 4 tail bits, that is n and since the rate, of course the number of coded bits that we will get is

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- These two encoders have the same set of codewords, but a different mapping between information sequence and codeword. N
- Now consider termination the codes by appending four bits to an information block of length $N - 4$ to return the encoder to all zero state. $N-4+4 = N$ $2N$

2 n.

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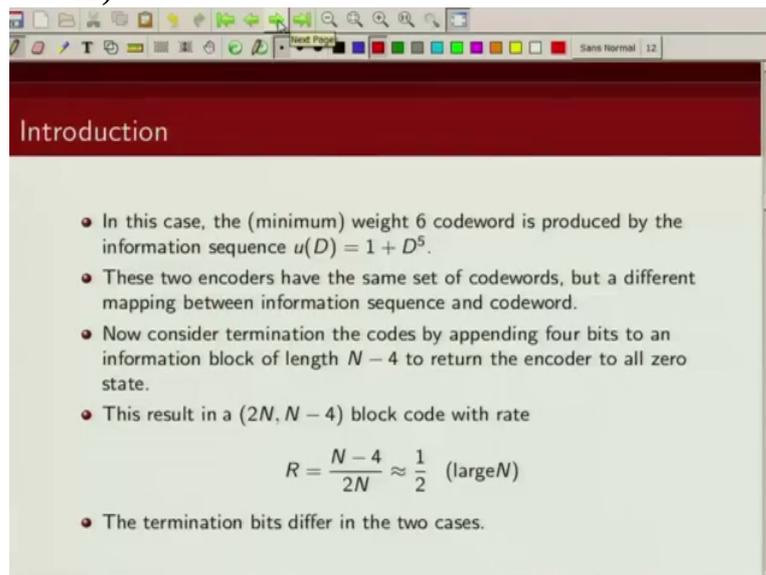
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- These two encoders have the same set of codewords, but a different mapping between information sequence and codeword.
- Now consider termination the codes by appending four bits to an information block of length $N - 4$ to return the encoder to all zero state.
- This result in a $(2N, N - 4)$ block code with rate

$$R = \frac{N - 4}{2N} \approx \frac{1}{2} \quad (\text{large } N)$$

So this will result in a $2n - 4$ block code. So this is roughly a rate half code if n is very large.

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Introduction

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$$R = \frac{N - 4}{2N} \approx \frac{1}{2} \quad (\text{large } N)$$

- The termination bits differ in the two cases.

Now please note the termination bit depends on the type of the encoder that we are using. So the termination

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bit that we will use for the feedback encoders are different from the termination bits that we will be using for a

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Introduction

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- The termination bits differ in the two cases.

feed forward encoder.

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Introduction

Weight	Multiplicity	Weight	Multiplicity
0	1	17	520
1	0	18	491
2	0	19	346
3	0	20	212
4	0	21	132
5	0	22	68
6	11	23	38
7	12	24	11
8	23	25	2
9	38	26	0
10	61	27	0
11	126	28	0
12	200	29	0
13	332	30	0
14	425	31	0
15	502	32	0
16	545	-	-

Table: Weight spectra of (32,12) terminated convolutional code

So in this I have drawn, so , consider m equal to 16. So this is

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Weight	Multiplicity	Weight	Multiplicity
0	1	17	520
1	0	18	491
2	0	19	346
3	0	20	212
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16	545	-	-

N=16

Table: Weight spectra of (32,12) terminated convolutional code

a 12 32, so n is 32 and k is 12 here. Remember we are using a generator matrix of memory order 4 so we are terminating it, so 4 (0) 4 tail bits. So this is a 32 12 terminated convolution code. And look at the weight spectrum of that. So there is a codeword all zero codeword. Then we have 11 codewords of weight 6, 12 codewords of weight 7, 23 codewords of weight 8, 38 codewords of weight 9 and so on.

Now let us compare

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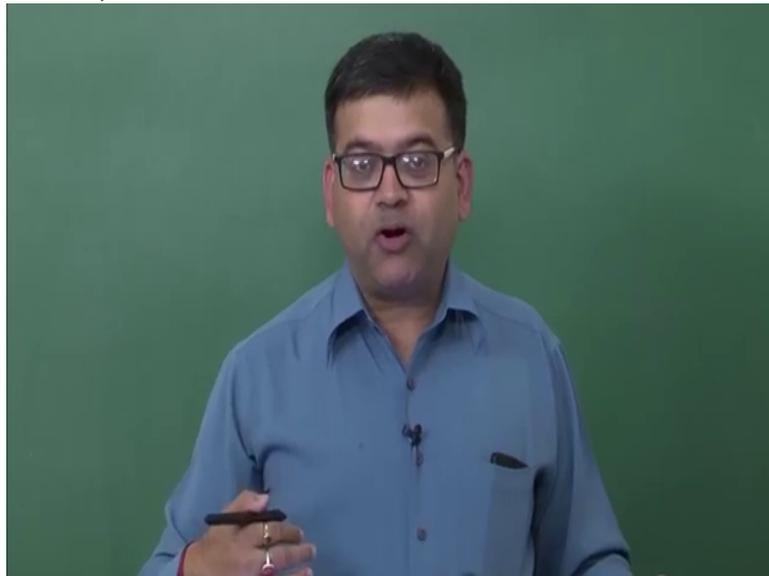
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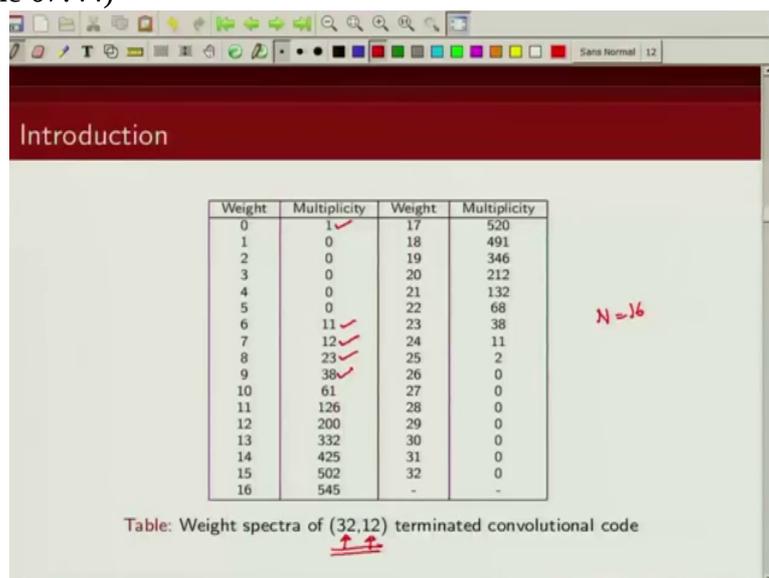
the weight distribution

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of a terminated convolutional code with a corresponding parallel concatenated code

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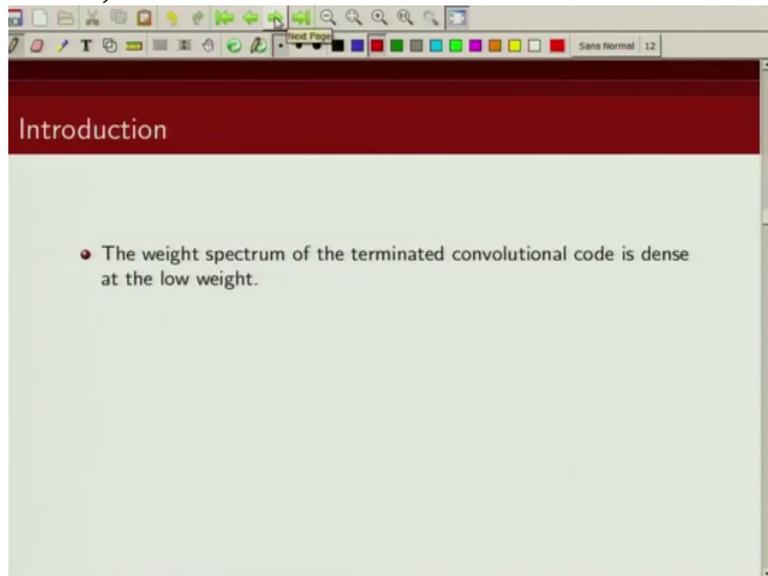


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Table: Weight spectra of (32,12) terminated convolutional code

with similar parameters. So

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the thing to be noted here is if you look here,

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Introduction

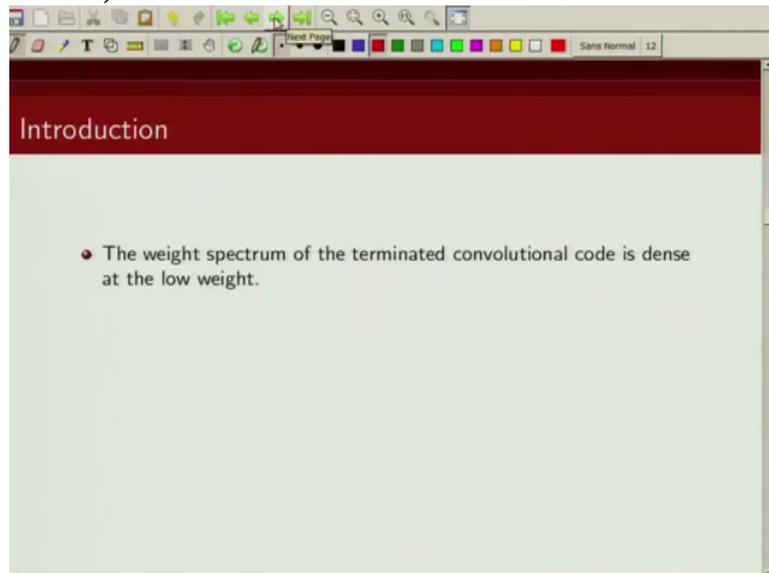
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$N=16$

Table: Weight spectra of (32,12) terminated convolutional code

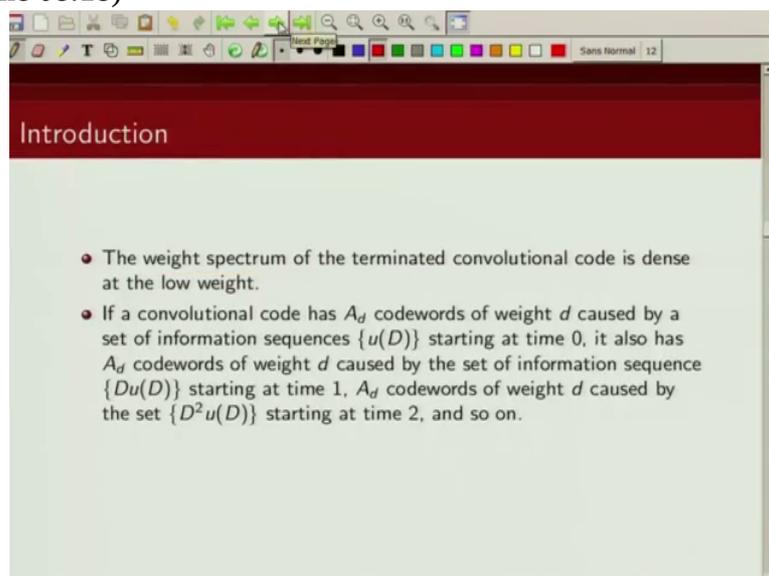
the multiplicity of codewords having small weight is quite high. There are 11 codewords of weight 6, 12 codewords of weight 7, 23 codewords of weight 8, 38 codewords of weight 9. So

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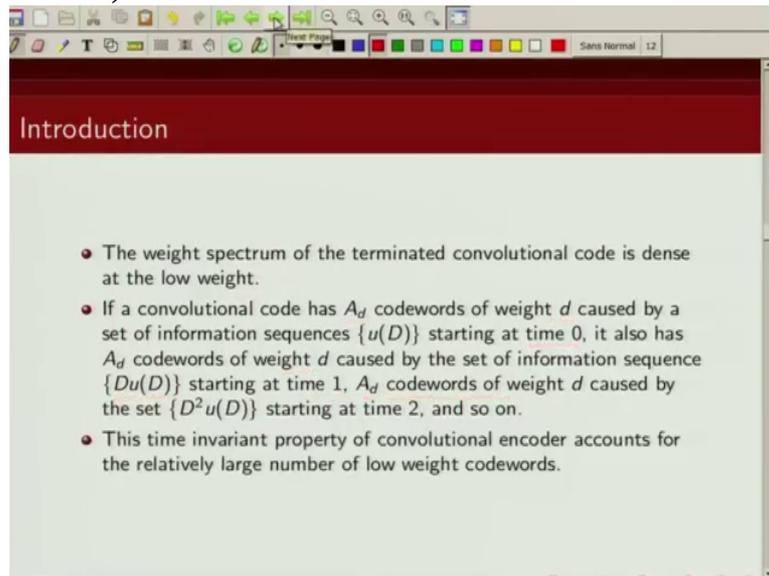
the statement I am making is spectrum of terminated convolutional code is dense at lower weight.

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Now if you have convolutional code which has 80 codewords of weight d caused by set of information sequence starting at time zero, then it also has these 80 codewords of weight d caused by set of information sequence of form $D u d$ at time 1 or 80 codewords of weight d caused by the information sequence of the form $D^2 u d$ starting at time 1 and so on.

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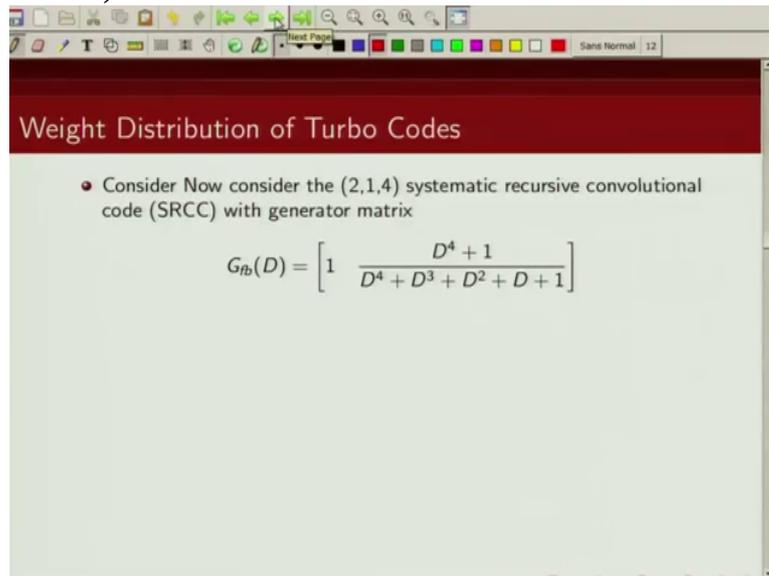
So this time invariant property of convolutional encoder accounts for their relatively

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large number of low weight codewords in the weight spectrum of the convolutional code.

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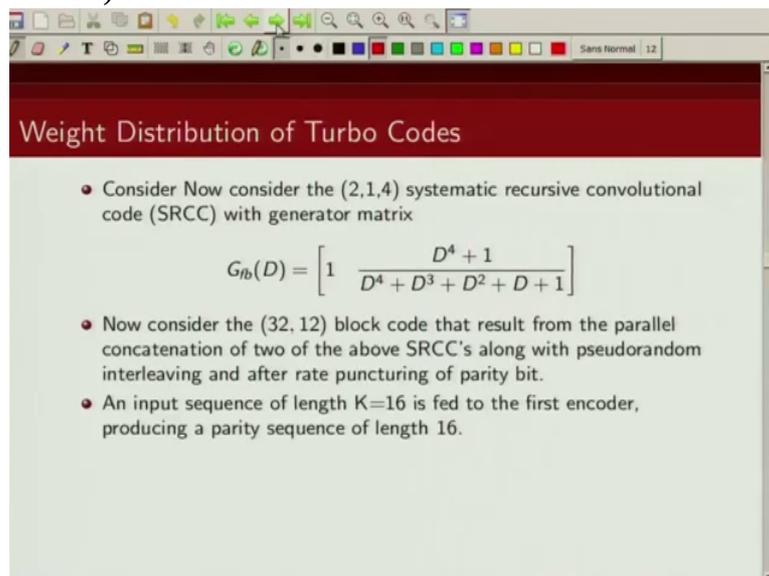
The slide is titled "Weight Distribution of Turbo Codes" and contains the following text:

- Consider Now consider the (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_{fb}(D) = \left[1 \quad \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \right]$$

Now, let us now consider the equivalent systematic feedback encoder, this is equivalent to the feed forward encoder generator matrix that we saw earlier. Now

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- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit.
- An input sequence of length $K=16$ is fed to the first encoder, producing a parity sequence of length 16.

we want to construct a parallel concatenation of 2 encoders which has the same parameters as the terminated convolutional code that we had. So how do we get it? So if you recall, turbo code is of the form like this. So this is encoder 1, this is encoder 2, and then this is your interleaver and then your parity stream coming encoder 1, parity stream coming encoder 2 and then you have these systematic bits

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Next to the equation is a block diagram of a feedback convolutional encoder. It shows an input bit entering a feedback loop that passes through a delay element (D) and a summing junction. The output of the summing junction goes through a feedback loop that passes through a delay element (D) and a summing junction. The output of the summing junction goes through a feedback loop that passes through a delay element (D) and a summing junction. The output of the summing junction goes through a feedback loop that passes through a delay element (D) and a summing junction. The output of the summing junction goes through a feedback loop that passes through a delay element (D) and a summing junction.

- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit.
- An input sequence of length K=16 is fed to the first encoder, producing a parity sequence of length 16.

right. So if you have k information bits coming in, here you are getting k information bits, here you are getting k information bits, here you are getting n minus k parity bits, here you are getting n minus k parity bits.

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- Consider Now consider the (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_{fb}(D) = \left[1 \quad \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \right]$$

Next to the equation is a block diagram of a feedback convolutional encoder. It shows an input bit entering a feedback loop that passes through a delay element (D) and a summing junction. The output of the summing junction goes through a feedback loop that passes through a delay element (D) and a summing junction. The output of the summing junction goes through a feedback loop that passes through a delay element (D) and a summing junction. The output of the summing junction goes through a feedback loop that passes through a delay element (D) and a summing junction. The output of the summing junction goes through a feedback loop that passes through a delay element (D) and a summing junction.

- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit.
- An input sequence of length K=16 is fed to the first encoder, producing a parity sequence of length 16.

Now let us consider a input of length 16 where 12 of them will be information bits Ok and 4 of them will be tail bits which we use

(Refer Slide Time 10:34)

The slide is titled "Weight Distribution of Turbo Codes" and contains the following text:

- Consider Now consider the (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_{rb}(D) = \begin{bmatrix} 1 & \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \end{bmatrix}$$

Next to the equation is a block diagram of a convolutional encoder with two outputs labeled $n-k$ and $n-k$.

- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit. *12 information*
- An input sequence of length K=16 is fed to the first encoder, producing a parity sequence of length 16. *4 tail bits*

to terminate this convolutional encoder. So we are considering input of size 16, 12 of them is information sequence and 4 of these are tail bits which are used to terminate this convolutional encoder. So this we are feeding our encoder 1, which is this encoder e 1,

(Refer Slide Time 10:55)

This slide is identical to the one above, showing the generator matrix and the description of the (32, 12) block code formed by concatenating two SRCCs.

Ok. Now input sequence, and we are feeding an input sequence and remember this is our rate one half code, this is rate one half code. If we feed the sequence of input bit 16, the parity bit that will come out will have parity bit of length 16. Because here k is half of n. So if we feed in information sequence of length 16, here we are getting

(Refer Slide Time 11:30)

Weight Distribution of Turbo Codes

- Consider Now consider the (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_{rb}(D) = \begin{bmatrix} 1 & \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \end{bmatrix}$$

Diagram: A block diagram of a (2,1,4) SRCC encoder. It shows an input sequence of length K entering a feedback loop with a delay element D. The loop contains a summing junction and a feedback path. The output sequence has length n-k.

- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit. *12 information*
- An input sequence of length K=16 is fed to the first encoder, producing a parity sequence of length 16. *4 tail bits*

also a parity sequence of length 16, because it is a rate half code. Next what's happening

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Weight Distribution of Turbo Codes

- Consider Now consider the (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_{rb}(D) = \begin{bmatrix} 1 & \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \end{bmatrix}$$

- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit.
- An input sequence of length K=16 is fed to the first encoder, producing a parity sequence of length 16.
- The interleaving pattern is given by

$$\pi_{16} = [0, 8, 15, 9, 4, 7, 11, 5, 1, 3, 14, 6, 13, 12, 10, 2]$$

$$(u'_0 = u_0, u'_1 = u_8, u'_2 = u_{15}, \dots, u'_{15} = u_2)$$

is, so this information sequence is interleaved

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Weight Distribution of Turbo Codes

- Consider Now consider the (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_b(D) = \begin{bmatrix} 1 & \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \end{bmatrix}$$

- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit. *12 information*
- An input sequence of length $K=16$ is fed to the first encoder, producing a parity sequence of length 16. *4 tail bits*

and fed to second encoder. So what the second encoder does, it is an interleaved version. Interleaved is nothing but it is a permuted version of this information sequence. So this is a 16 sequence that comes here and this also generates a parity sequence of length 16.

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Weight Distribution of Turbo Codes

- Consider Now consider the (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_b(D) = \begin{bmatrix} 1 & \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \end{bmatrix}$$

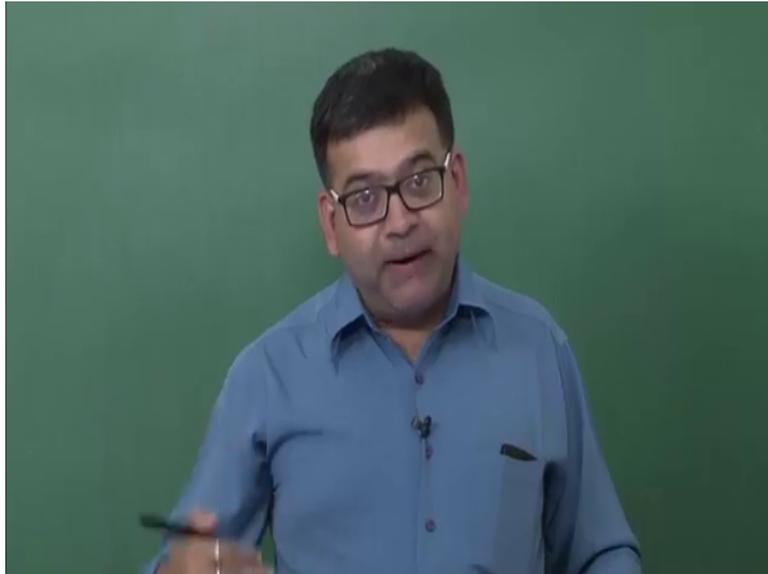
- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit.
- An input sequence of length $K=16$ is fed to the first encoder, producing a parity sequence of length 16.
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$$\pi_{16} = [0, 8, 15, 9, 4, 7, 11, 5, 1, 3, 14, 6, 13, 12, 10, 2]$$

$$(u'_0 = u_0, u'_1 = u_8, u'_2 = u_{15}, \dots, u'_{15} = u_2)$$

So we need to specify, remember we are trying to get to the distance profile of a parallel concatenated code and the distance profile of parallel concatenated code

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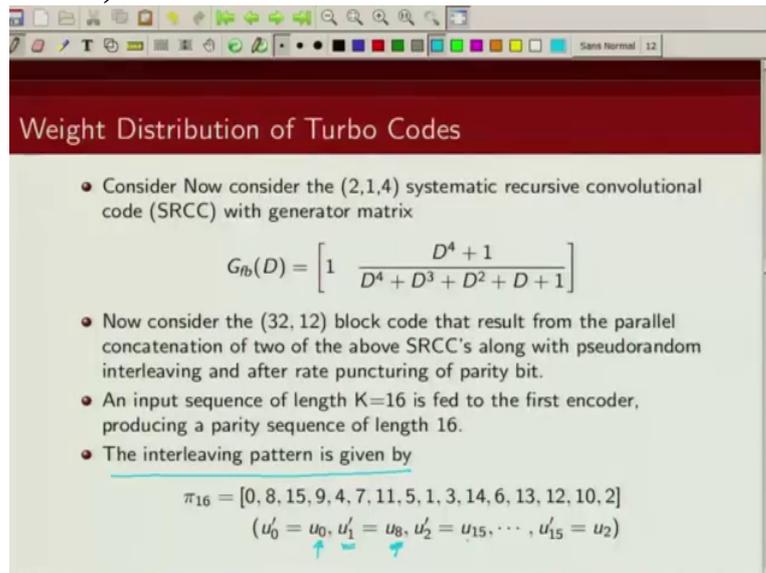
depends on what type of interleaver we are using. In this particular example

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A screenshot of a presentation slide. The title is "Weight Distribution of Turbo Codes" in white text on a dark red background. The slide content is on a light green background. It contains a list of bullet points and mathematical expressions. The first bullet point mentions a (2,1,4) SRCC with a generator matrix. The second bullet point mentions a (32, 12) block code. The third bullet point mentions an input sequence of length K=16. The fourth bullet point mentions an interleaving pattern. The slide is displayed in a window with a standard toolbar and a "Next Page" button.

we are specifying one such interleaver pattern. So you can see here I have specified here so 0 layer means the bit stays here in the interleaved version the first bit belongs, the second bit belongs to the input u_8 and like that. So u_0 s are the original input sequence

(Refer Slide Time 12:53)



The slide is titled "Weight Distribution of Turbo Codes" and contains the following content:

- Consider Now consider the (2,1,4) systematic recursive convolutional code (SRCC) with generator matrix

$$G_b(D) = \left[1 \quad \frac{D^4 + 1}{D^4 + D^3 + D^2 + D + 1} \right]$$

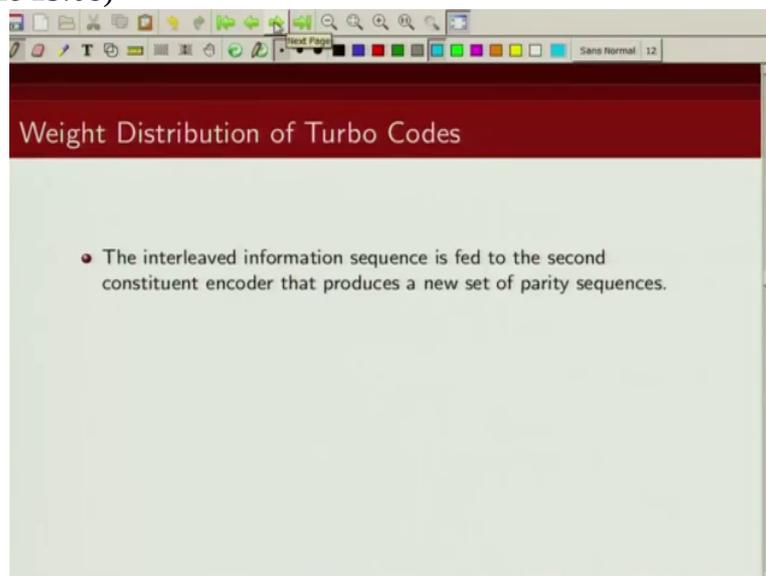
- Now consider the (32, 12) block code that result from the parallel concatenation of two of the above SRCC's along with pseudorandom interleaving and after rate puncturing of parity bit.
- An input sequence of length $K=16$ is fed to the first encoder, producing a parity sequence of length 16.
- The interleaving pattern is given by

$$\pi_{16} = [0, 8, 15, 9, 4, 7, 11, 5, 1, 3, 14, 6, 13, 12, 10, 2]$$
$$(u'_0 = u_0, u'_1 = u_8, u'_2 = u_{15}, \dots, u'_{15} = u_2)$$

Blue arrows point from the u'_i terms to the corresponding u_j terms in the sequence below.

and u_i 's are the interleaved information sequence. So this is the interleaving pattern I am using for this particular example to derive the weight distribution of this parallel concatenated code.

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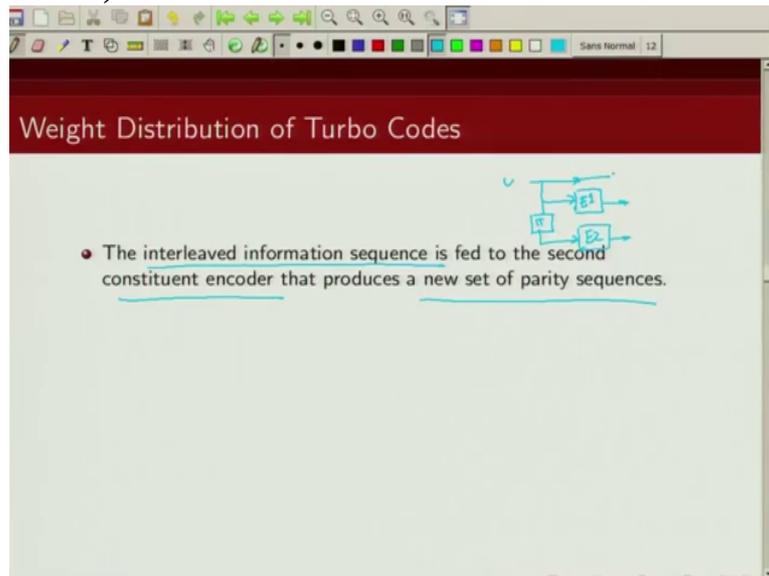


The slide is titled "Weight Distribution of Turbo Codes" and contains the following content:

- The interleaved information sequence is fed to the second constituent encoder that produces a new set of parity sequences.

Now this interleaved information sequence is fed to the second constituent encoder and which produces a new set of parity sequences. Now as you know the parallel concatenating encoder we drew is actually a example of a rate one third code. This was our parallel concatenated code. Our objective is to first compare the weight spectrum

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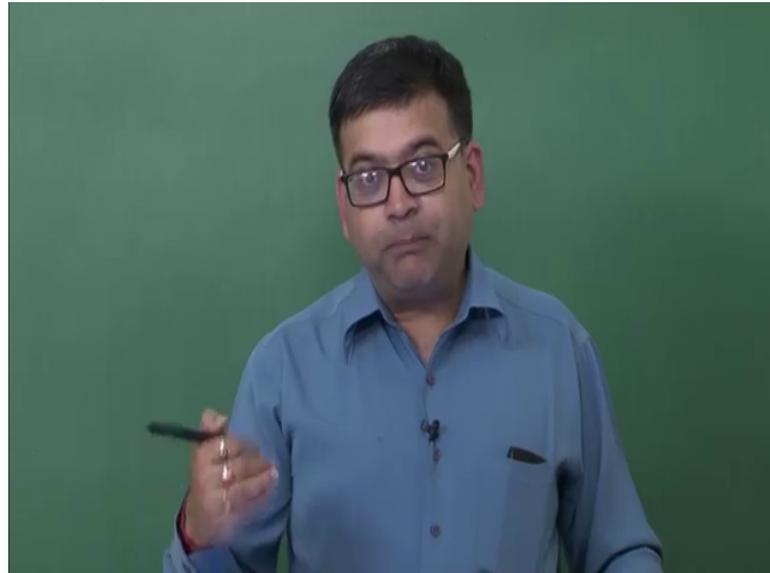
Weight Distribution of Turbo Codes

- The interleaved information sequence is fed to the second constituent encoder that produces a new set of parity sequences.

The slide also features a small block diagram with two rectangular blocks labeled E1 and E2. An input arrow enters E1 from the left, and an arrow connects E1 to E2. An output arrow exits E2 to the right. There is a checkmark above the diagram.

of a convolutional code to that of a turbo code. And the convolutional code that you studied was a rate half code. And this is a rate one third code. So what do we need to do to make it a rate one half? So we are going to puncture

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the parity bits. So we are going to send alternatively parity bits from

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Weight Distribution of Turbo Codes

- The interleaved information sequence is fed to the second constituent encoder that produces a new set of parity sequences.

output of the first encoder or output of the second encoder. So for example, let's call this v_0 , v_1 and v_2 . So what we are

(Refer Slide Time 14:15)

Weight Distribution of Turbo Codes

- The interleaved information sequence is fed to the second constituent encoder that produces a new set of parity sequences.

going to do is at a particular time instance we transmit v_0 and v_1 and would not transmit this parity bit

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Weight Distribution of Turbo Codes

- The interleaved information sequence is fed to the second constituent encoder that produces a new set of parity sequences.

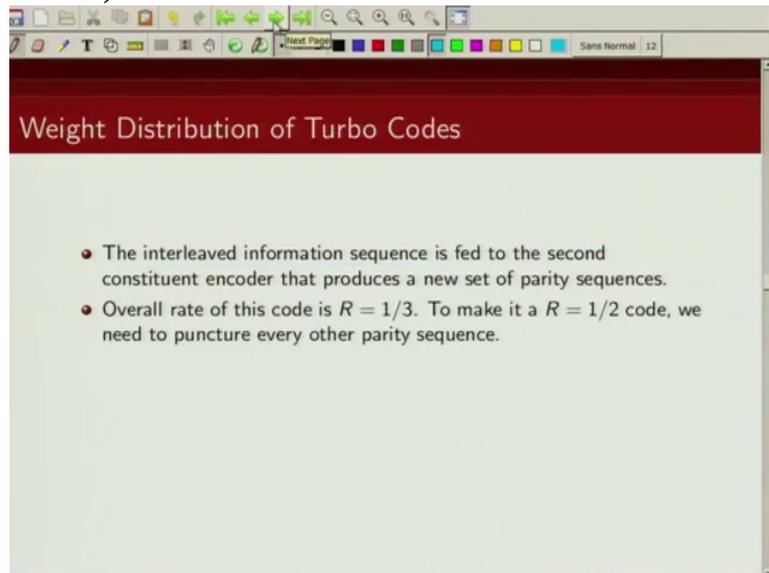
and the second instance will transmit this information sequence, we would not send this parity bit, we will send parity bit from this encoder.

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Weight Distribution of Turbo Codes

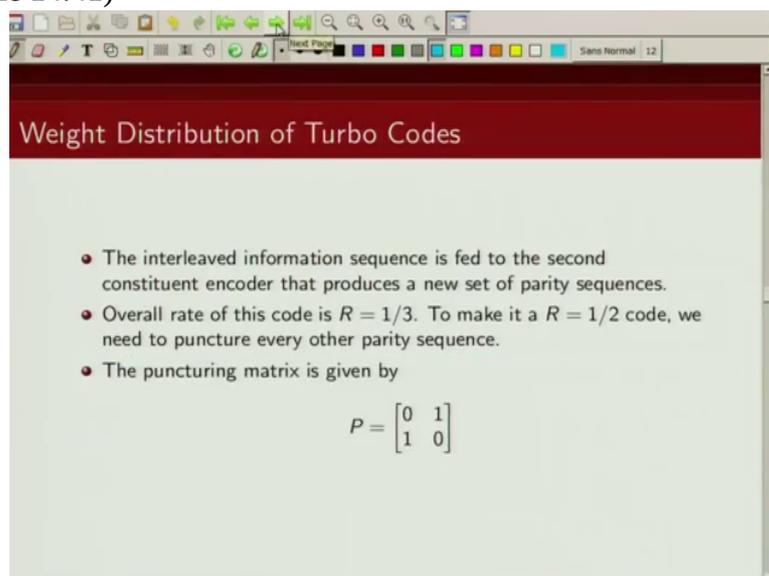
- The interleaved information sequence is fed to the second constituent encoder that produces a new set of parity sequences.

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So we need to puncture every other parity bit in each of these streams to make it a rate one half

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and we denote a puncturing pattern like this. So you can view

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The slide is titled "Weight Distribution of Turbo Codes" and contains the following text:

- The interleaved information sequence is fed to the second constituent encoder that produces a new set of parity sequences.
- Overall rate of this code is $R = 1/3$. To make it a $R = 1/2$ code, we need to puncture every other parity sequence.
- The puncturing matrix is given by

$$P = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

and we denote a puncturing pattern like this. So you can view this, if there is a zero, we are not transmitting that bit. If there is a 1, we are transmitting this bit. So you can look at, view look at, view it like this. At odd time instances I am sending parity bit from the second encoder where as in the even time instances I am sending

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The slide is titled "Weight Distribution of Turbo Codes" and contains the following text:

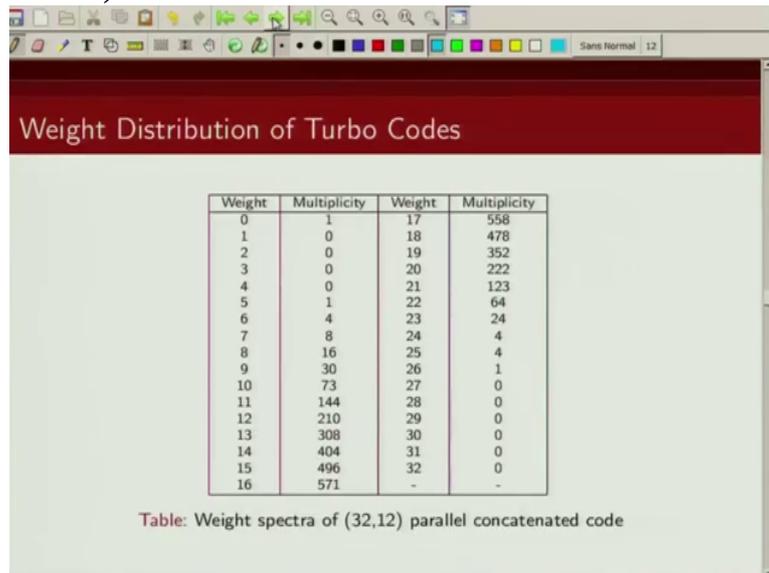
- The interleaved information sequence is fed to the second constituent encoder that produces a new set of parity sequences.
- Overall rate of this code is $R = 1/3$. To make it a $R = 1/2$ code, we need to puncture every other parity sequence.
- The puncturing matrix is given by

$$P = \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}$$

In this version of the slide, blue arrows point to the '1' in the top-right cell and the '1' in the bottom-left cell of the matrix.

parity bit from the first encoder. So this will make the overall rate of this parallel concatenated code as rate half which is same as the original convolutional code that

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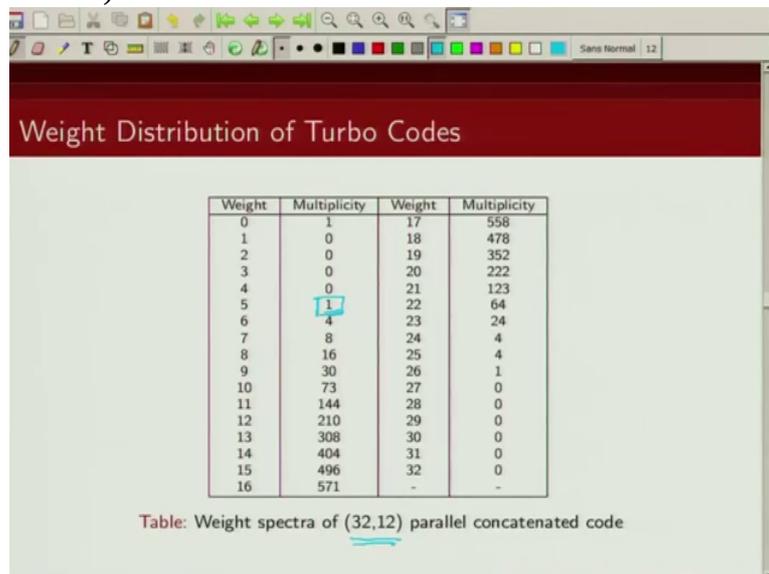


Weight	Multiplicity	Weight	Multiplicity
0	1	17	558
1	0	18	478
2	0	19	352
3	0	20	222
4	0	21	123
5	1	22	64
6	4	23	24
7	8	24	4
8	16	25	4
9	30	26	1
10	73	27	0
11	144	28	0
12	210	29	0
13	308	30	0
14	404	31	0
15	496	32	0
16	571	-	-

Table: Weight spectra of (32,12) parallel concatenated code

we were talking about. Now look at the way the distribution of 32 12 parallel concatenated code which we just talked about. Now note in case of convolutional code the minimum d free was 6. But here we see there is a codeword

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Weight	Multiplicity	Weight	Multiplicity
0	1	17	558
1	0	18	478
2	0	19	352
3	0	20	222
4	0	21	123
5	1	22	64
6	4	23	24
7	8	24	4
8	16	25	4
9	30	26	1
10	73	27	0
11	144	28	0
12	210	29	0
13	308	30	0
14	404	31	0
15	496	32	0
16	571	-	-

Table: Weight spectra of (32,12) parallel concatenated code

of weight 5. However the point to be noted is this. Now let's look at the multiplicity of codewords of low weight. In case of the example we considered earlier of the terminated convolutional code, in that case we had 11 codewords of

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Weight	Multiplicity	Weight	Multiplicity
0	1	17	558
1	0	18	478
2	0	19	352
3	0	20	222
4	0	21	123
5	1	22	64
6	4	23	24
7	8	24	4
8	16	25	4
9	30	26	1
10	73	27	0
11	144	28	0
12	210	29	0
13	308	30	0
14	404	31	0
15	496	32	0
16	571	-	-

Table: Weight spectra of (32,12) parallel concatenated code

weight 6 whereas here we are having only 4.

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Weight	Multiplicity	Weight	Multiplicity
0	1	17	558
1	0	18	478
2	0	19	352
3	0	20	222
4	0	21	123
5	1	22	64
6	4	23	24
7	8	24	4
8	16	25	4
9	30	26	1
10	73	27	0
11	144	28	0
12	210	29	0
13	308	30	0
14	404	31	0
15	496	32	0
16	571	-	-

Table: Weight spectra of (32,12) parallel concatenated code

We were having 12 codewords of weight 7, here we are having 8.

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Weight	Multiplicity	Weight	Multiplicity
0	1	17	558
1	0	18	478
2	0	19	352
3	0	20	222
4	0	21	123
5	1	22	64
6	4	23	24
7	8	24	4
8	16	25	4
9	30	26	1
10	73	27	0
11	144	28	0
12	210	29	0
13	308	30	0
14	404	31	0
15	496	32	0
16	571	-	-

Table: Weight spectra of (32,12) parallel concatenated code

We were having 23 codewords of weight 8, here we are having 16.

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Weight	Multiplicity	Weight	Multiplicity
0	1	17	558
1	0	18	478
2	0	19	352
3	0	20	222
4	0	21	123
5	1	22	64
6	4	23	24
7	8	24	4
8	16	25	4
9	30	26	1
10	73	27	0
11	144	28	0
12	210	29	0
13	308	30	0
14	404	31	0
15	496	32	0
16	571	-	-

Table: Weight spectra of (32,12) parallel concatenated code

We were having 38 codewords of weight 9, here we are having 30.

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Weight	Multiplicity	Weight	Multiplicity
0	1	17	558
1	0	18	478
2	0	19	352
3	0	20	222
4	0	21	123
5	1	22	64
6	4	23	24
7	8	24	4
8	16	25	4
9	30	26	1
10	73	27	0
11	144	28	0
12	210	29	0
13	308	30	0
14	404	31	0
15	496	32	0
16	571	-	-

Table: Weight spectra of $(32,12)$ parallel concatenated code

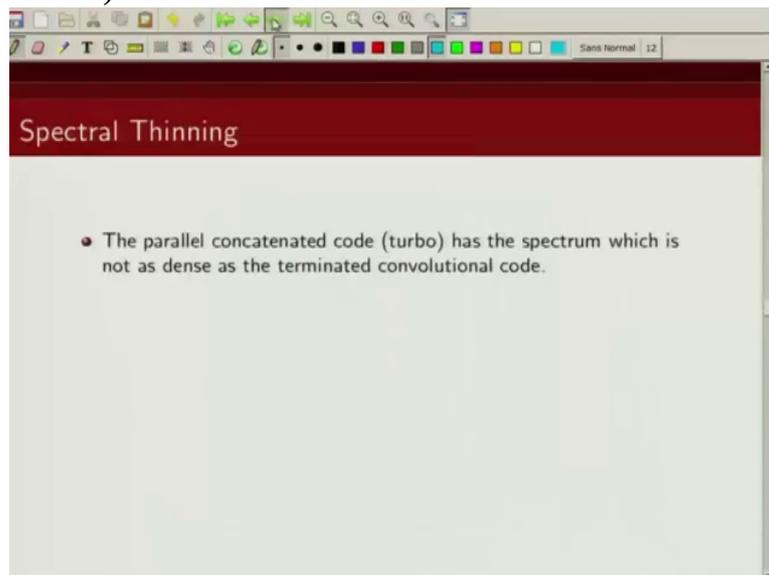
So note that in case of this concatenated code we are noticing that

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as a result of this interleaver the multiplicity of low weight codeword is reduced. Now we considered a very simple example with interleaver size of just 16. If we increase the interleaver size, this effect is even more pronounced. And this is

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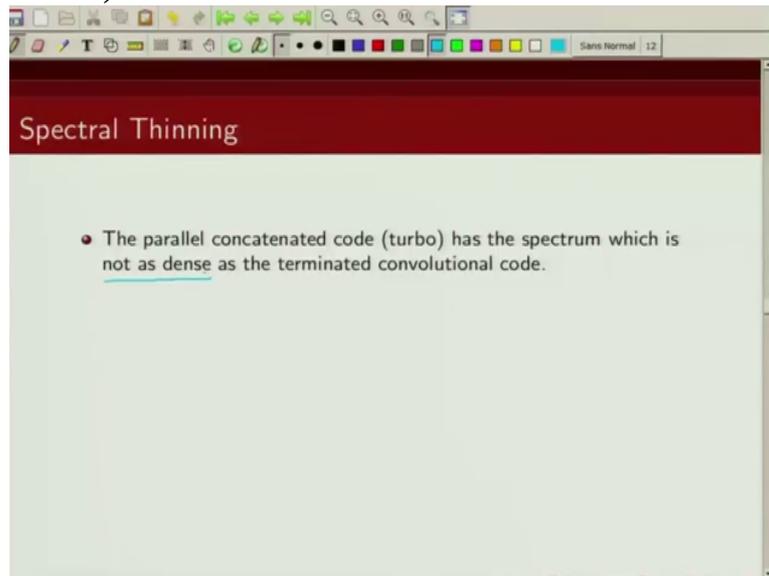
what we call spectral thinning. So parallel concatenate code has a spectrum which is not as dense as convolutional code. In fact we note, even with a very small example where we consider an interleaver size

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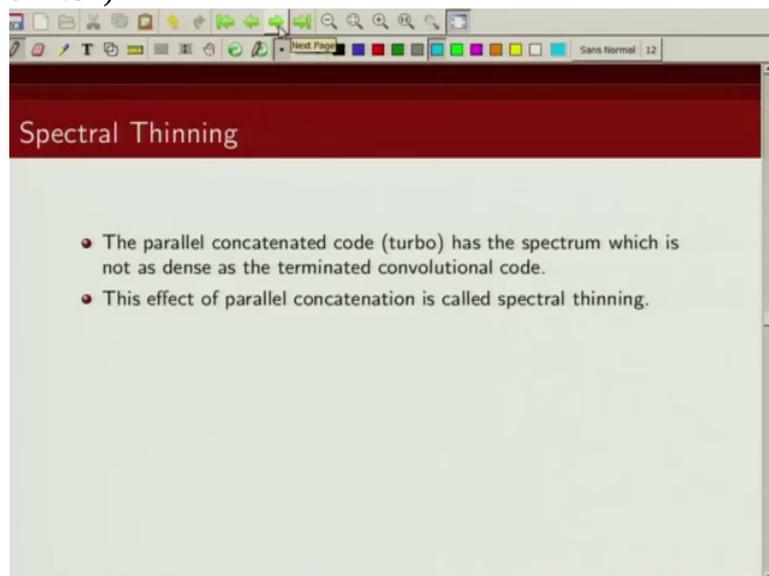
of 16, we saw that low weight codewords, the number of low weight codewords, the multiplicity of low weight codewords has reduced in case of

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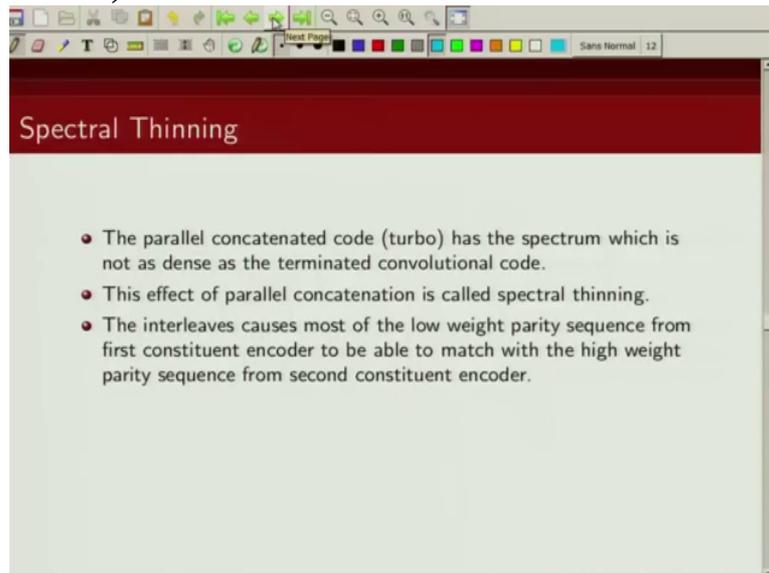
turbo codes.

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Now this effect of parallel concatenation where you know the multiplicity of low weight codewords getting reduced is known as spectral thinning. And this is observed when we are using systematic feedback encoder in convolutional code as constituent encoders.

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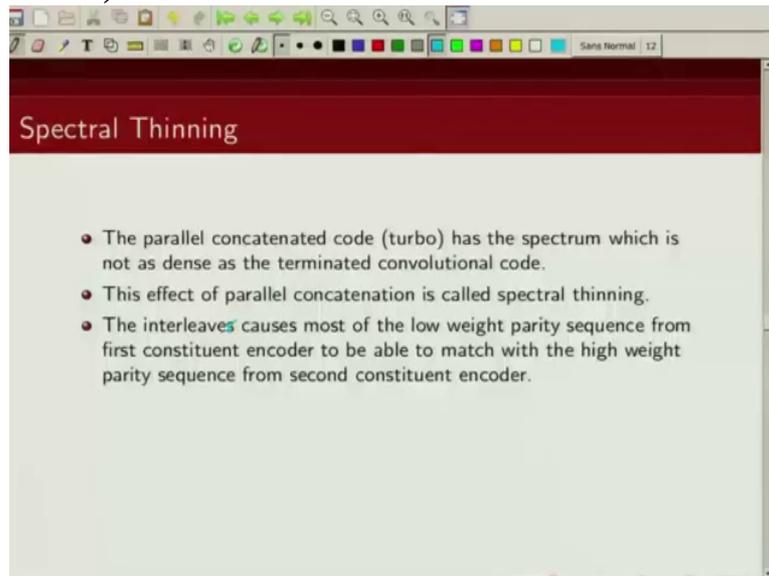
Now interleaver causes most the low weight parity sequence of the first encoder to be able to match to the high weight parity sequence from the second encoder. So the idea of design of interleaver

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is such that if the parity sequence coming out of the first encoder has small weight then the interleaved sequence when it is entering the second encoder it should produce a larger weight sequence so that overall code weight sequence is

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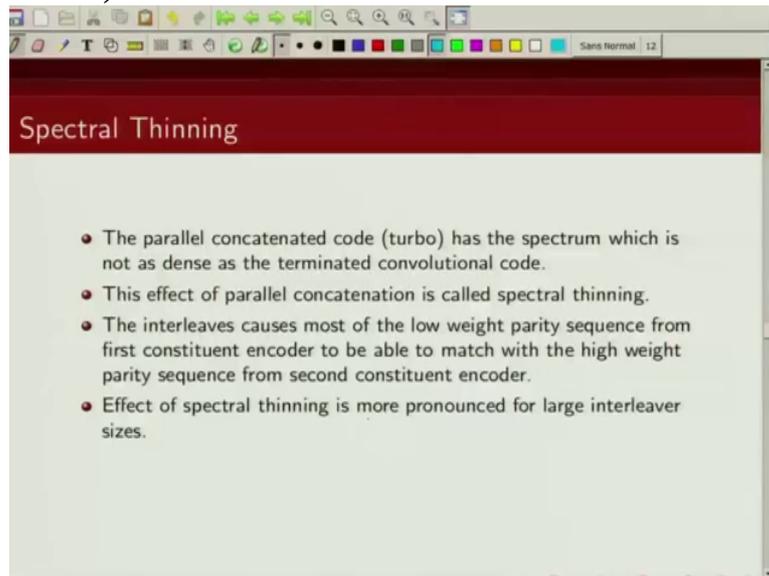
good. Now we just showed the effect of spectral thinning

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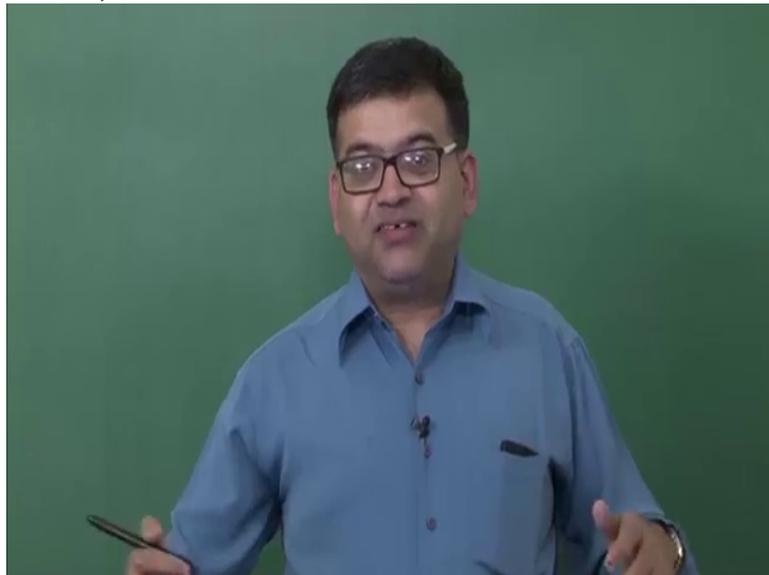
with the help of very small interleaver size of 16. Now this effect is more

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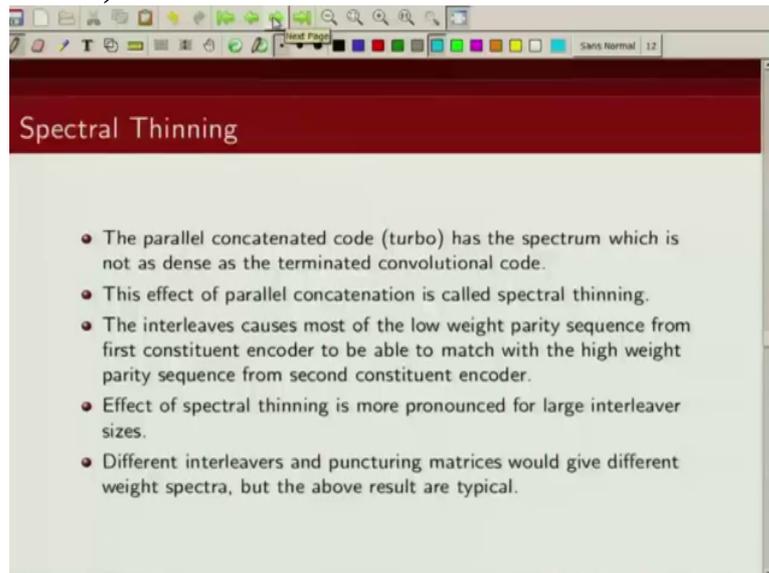
pronounced if we take larger interleaver. So let's say if I take interleaver size of 100

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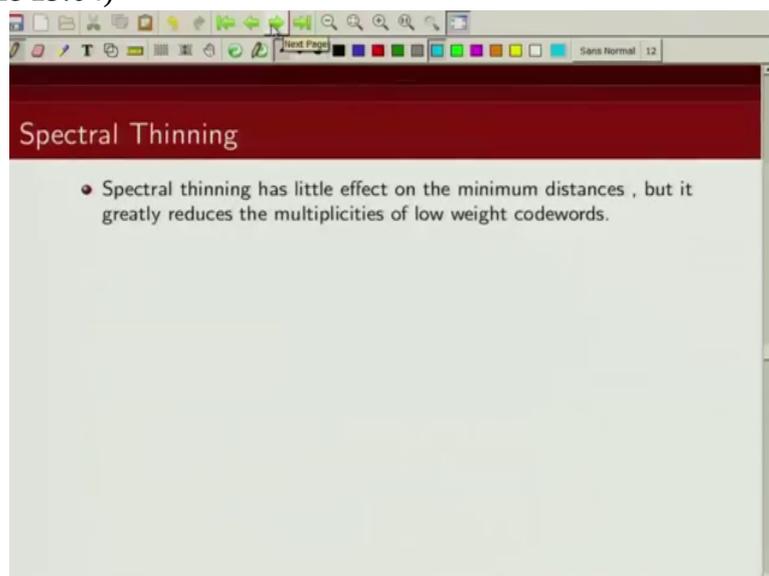
1000, 10000 you will see much pronounced effect of this spectral thinning. Now

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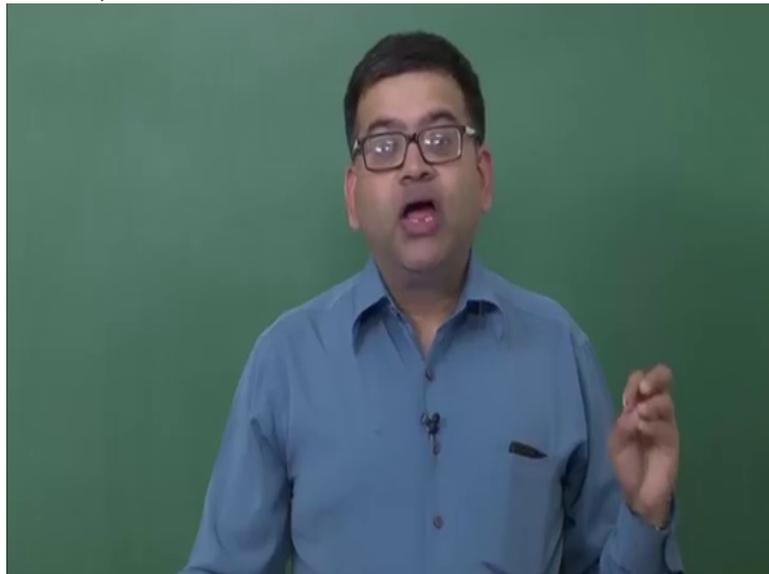
different interleavers and puncturing patterns would give different weight spectrum. But this result is very typical of turbo codes.

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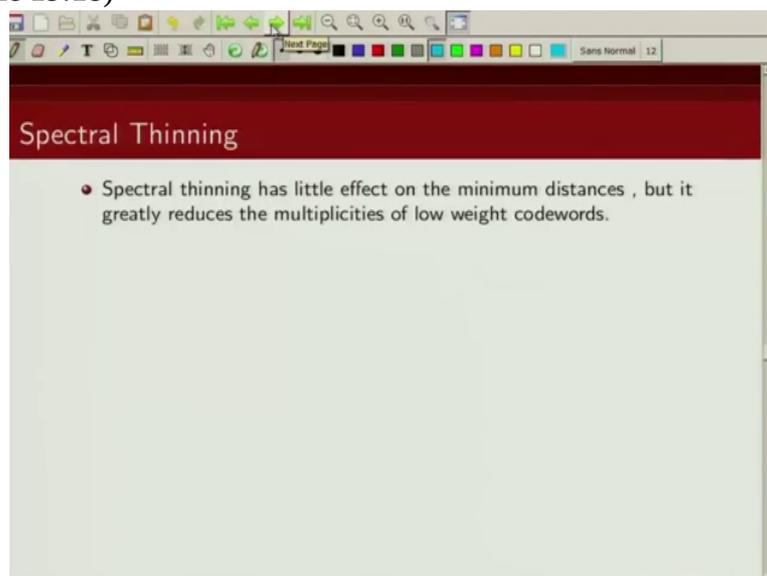
Now spectral thinning has very little effect on

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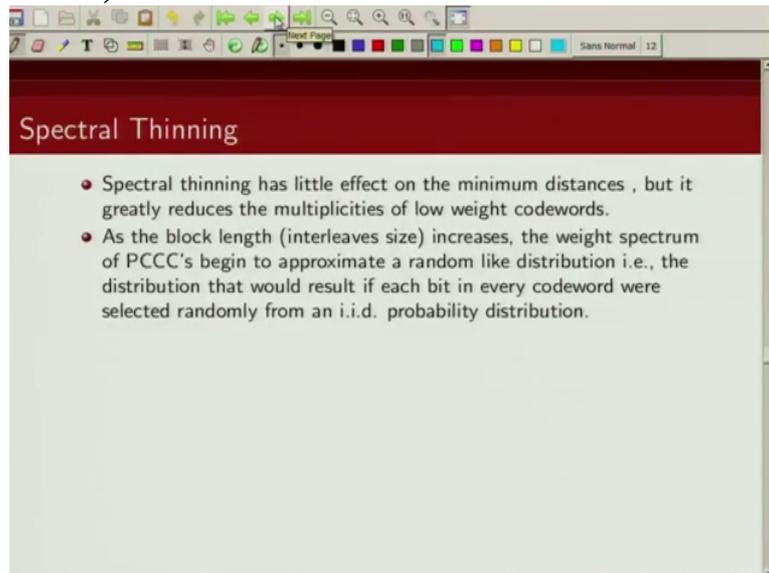
minimum distance of the turbo code. However it greatly reduces the multiplicity of low weight code words.

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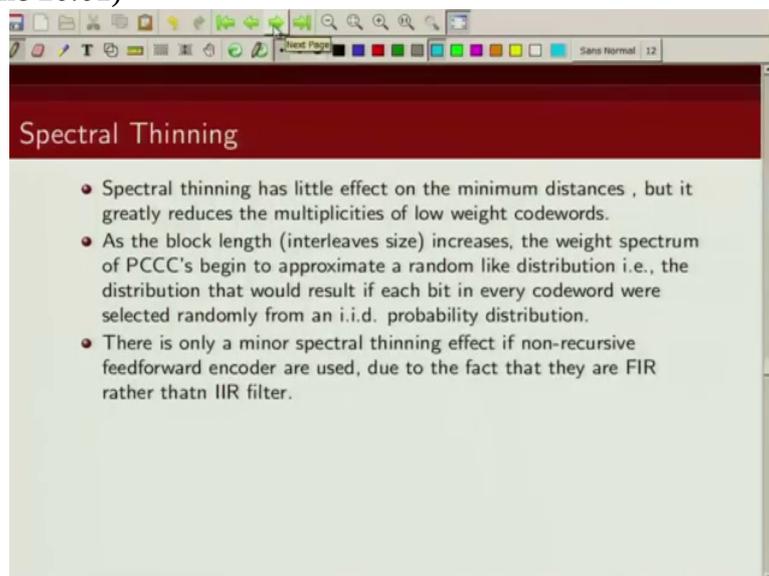
In the example that we considered, we saw that the minimum distance was 5 where as the convolutional code, for the convolutional code it was 6 but we only had one code word of weight 5 where as if you look at codewords, a number of codewords of weight 6, 7,8, 9 it was reduced in the case of parallel concatenated code.

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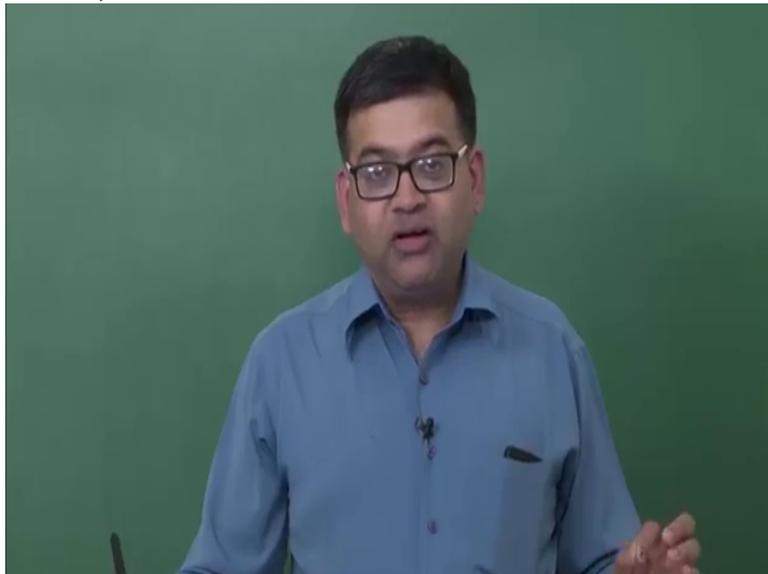
So as the block size increases, as the interleaver size increases the weight spectrum of turbo code looks more like random distribution so it is like the weight distribution as if every bit of the codeword was randomly selected according to some i i d distribution.

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Now the example that we took,

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we had considered the feedback systematic encoder. If you replace that feedback encoder with a feed forward encoder we see a very small spectral thinning effect even with large block sizes. So

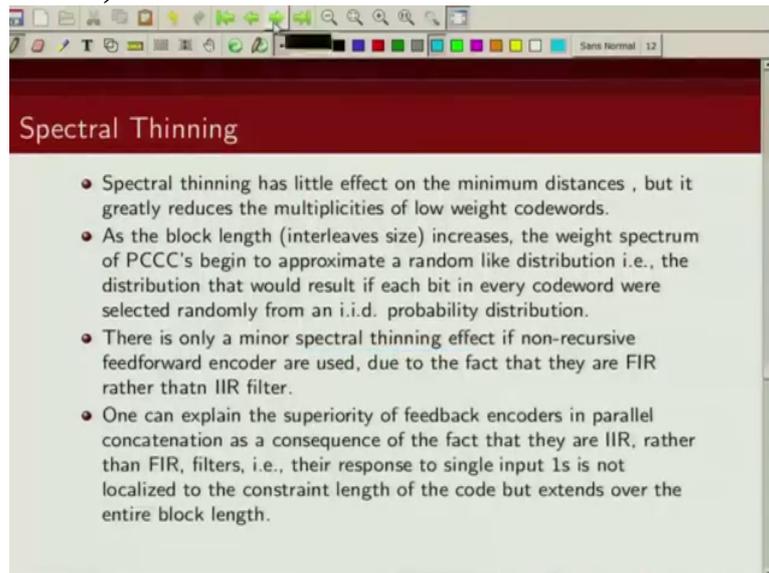
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A screenshot of a presentation slide. The slide has a dark red header with the title "Spectral Thinning" in white. The main content area is light green and contains three bullet points. The slide is shown within a software window that has a toolbar at the top with various icons and a status bar at the bottom that says "Next Page" and "Sans Normal 12".

- Spectral thinning has little effect on the minimum distances, but it greatly reduces the multiplicities of low weight codewords.
- As the block length (interleaves size) increases, the weight spectrum of PCCC's begin to approximate a random like distribution i.e., the distribution that would result if each bit in every codeword were selected randomly from an i.i.d. probability distribution.
- There is only a minor spectral thinning effect if non-recursive feedforward encoders are used, due to the fact that they are FIR rather than IIR filter.

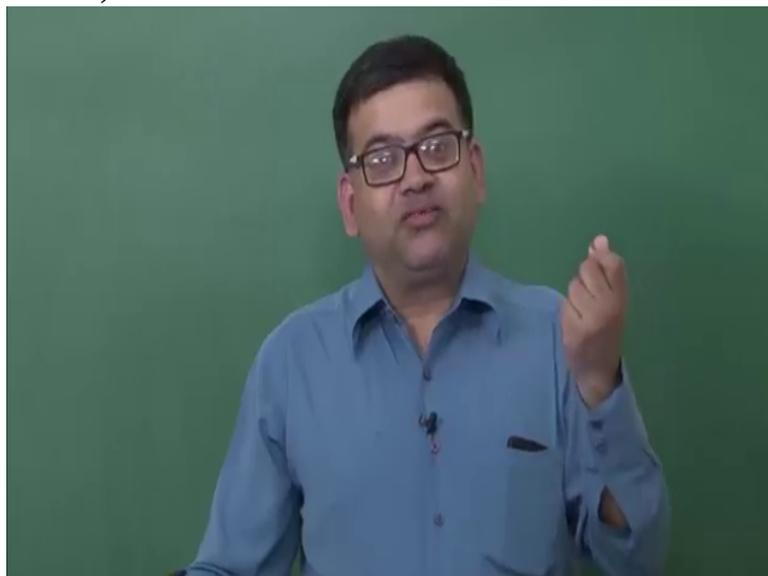
this spectral thinning effect is pronounced in case of a recursive encoder, not in a case of non recursive encoder.

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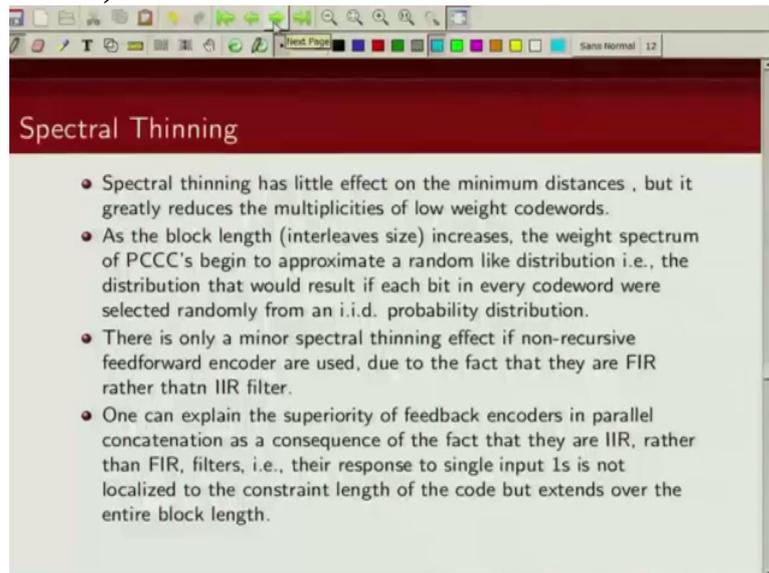
Now we can explain why feedback encoders are better if you feed in a way to one sequence to this

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feedback encoders they would not get terminated. So they will result in large weight output sequences whereas if we feed a weight 1s sequence to a feed forward encoder it can get terminated and we may get lower weight

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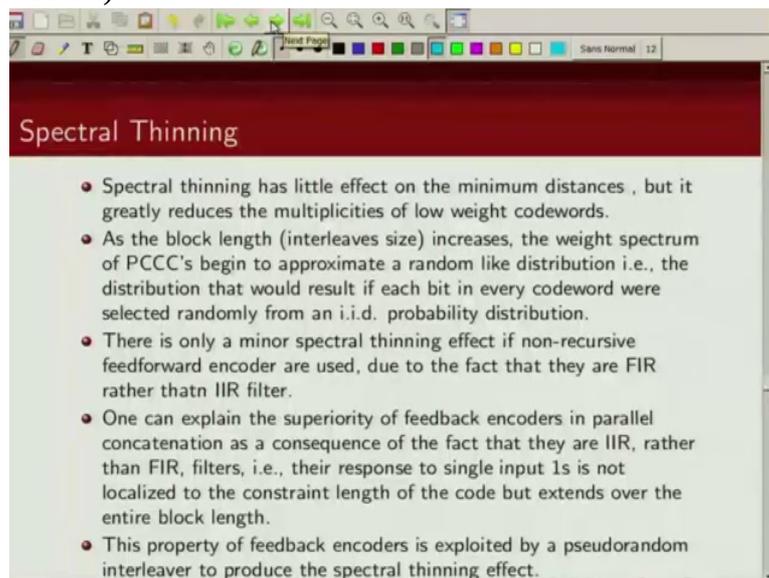


The screenshot shows a presentation slide with a dark red header containing the title "Spectral Thinning". Below the header, there are four bullet points in black text on a light grey background. The slide is displayed within a window that has a standard operating system taskbar at the top with various icons and a "Next Page" button.

- Spectral thinning has little effect on the minimum distances , but it greatly reduces the multiplicities of low weight codewords.
- As the block length (interleaves size) increases, the weight spectrum of PCCC's begin to approximate a random like distribution i.e., the distribution that would result if each bit in every codeword were selected randomly from an i.i.d. probability distribution.
- There is only a minor spectral thinning effect if non-recursive feedforward encoder are used, due to the fact that they are FIR rather than IIR filter.
- One can explain the superiority of feedback encoders in parallel concatenation as a consequence of the fact that they are IIR, rather than FIR, filters, i.e., their response to single input 1s is not localized to the constraint length of the code but extends over the entire block length.

sequences. So the response to single input 1 in case of a feed forward encoder is localized where as for a feedback encoder it extends over this entire block size. Now this

(Refer Slide Time 21:15)



This screenshot is identical to the one above, showing a presentation slide titled "Spectral Thinning". It contains five bullet points. The fifth bullet point, which was not present in the previous slide, states: "This property of feedback encoders is exploited by a pseudorandom interleaver to produce the spectral thinning effect."

- Spectral thinning has little effect on the minimum distances , but it greatly reduces the multiplicities of low weight codewords.
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- There is only a minor spectral thinning effect if non-recursive feedforward encoder are used, due to the fact that they are FIR rather than IIR filter.
- One can explain the superiority of feedback encoders in parallel concatenation as a consequence of the fact that they are IIR, rather than FIR, filters, i.e., their response to single input 1s is not localized to the constraint length of the code but extends over the entire block length.
- This property of feedback encoders is exploited by a pseudorandom interleaver to produce the spectral thinning effect.

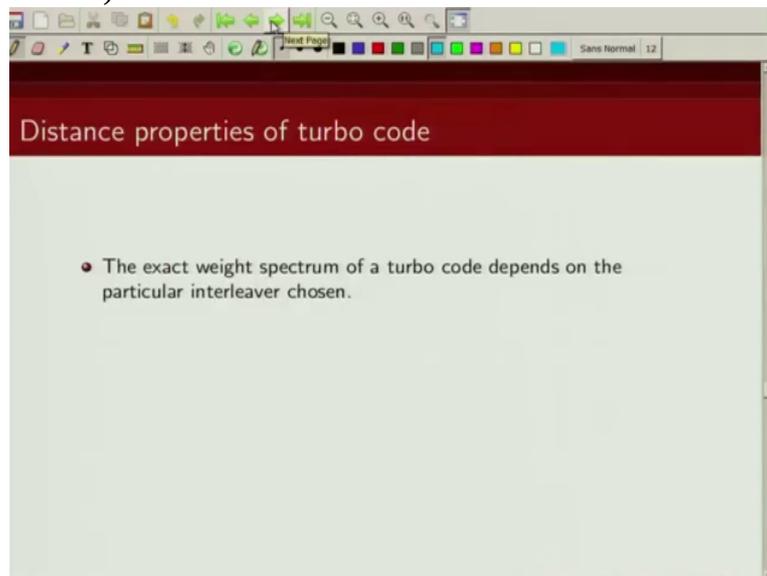
property of feedback encoders is exploited by the pseudo random interleaver to produce this spectral thinning effect.

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So what we have done so far is we have, with the help of the simple example for 32 12 terminated convolutional code and in 32 12 turbo codes we had shown the effect of spectral thinning. Now

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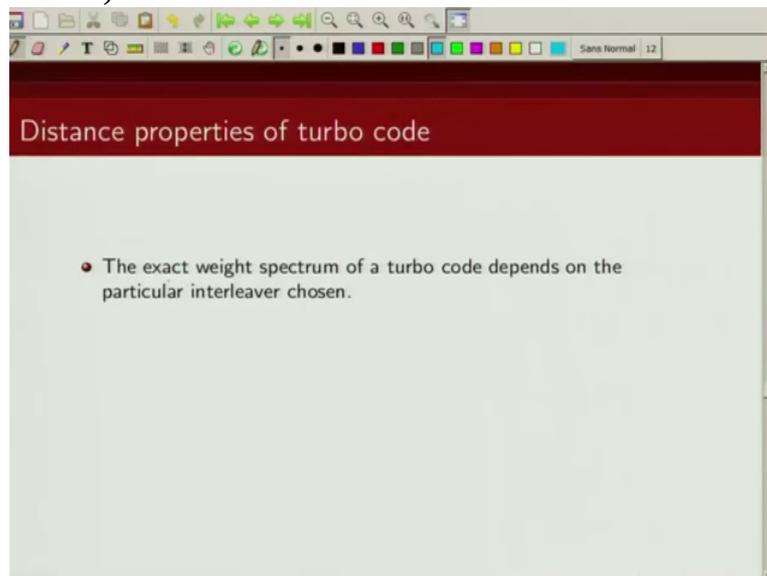
you want to find out what is the, how can we find out the weight distribution of a parallel concatenated code.

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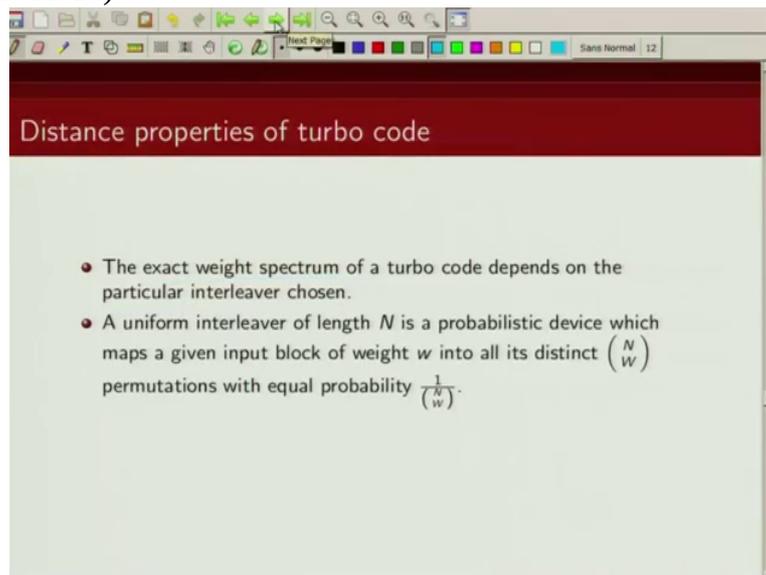
So again we will take a very simple example and with the help of this example we will try to illustrate how we can compute weight distribution of a parallel concatenated code. Now as you can understand

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the exact weight spectrum of a turbo code depends on a particular interleaver. So the weight spectrum is dependent on the interleaver we are using.

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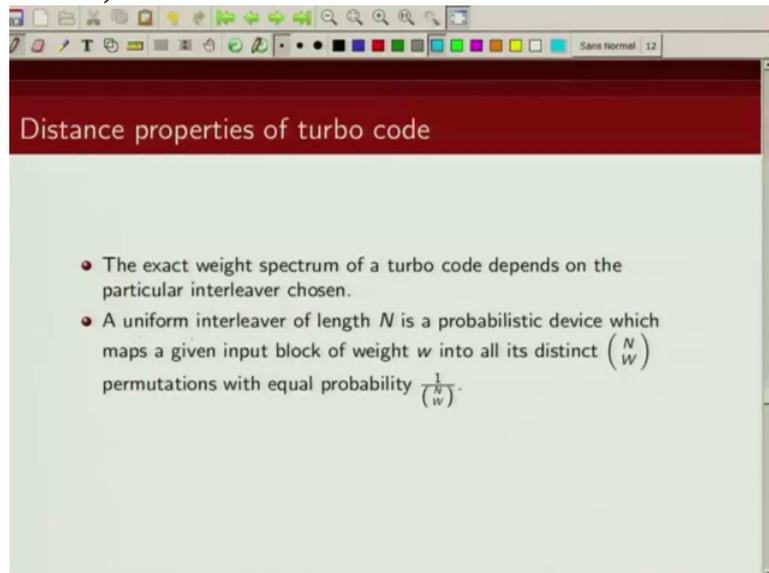
So we are going to make use of a device, we call it uniform interleaver. So this will help us

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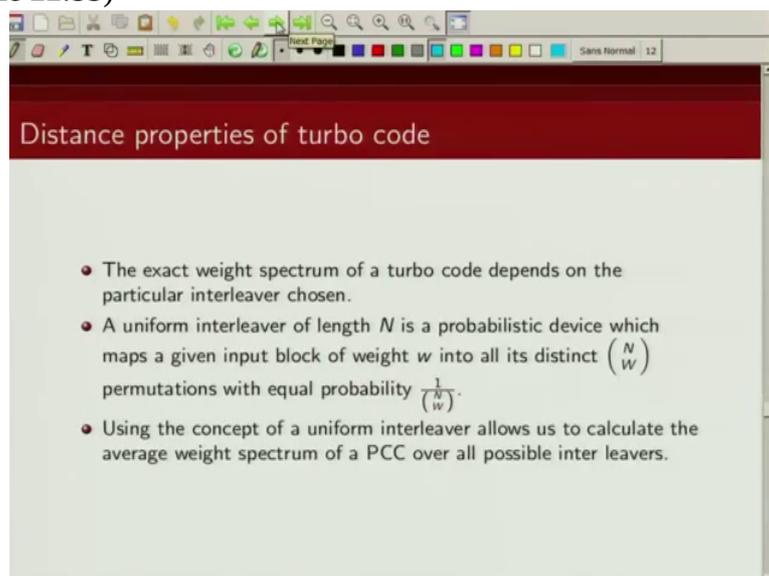
come up with the weight enumerating function of this parallel concatenated code averaged over all interleaver. So what is the uniform interleaver?

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A uniform interleaver of length n is a probabilistic device which maps an input block of weight w into all distinct $n w$ permutations with equal probability. So

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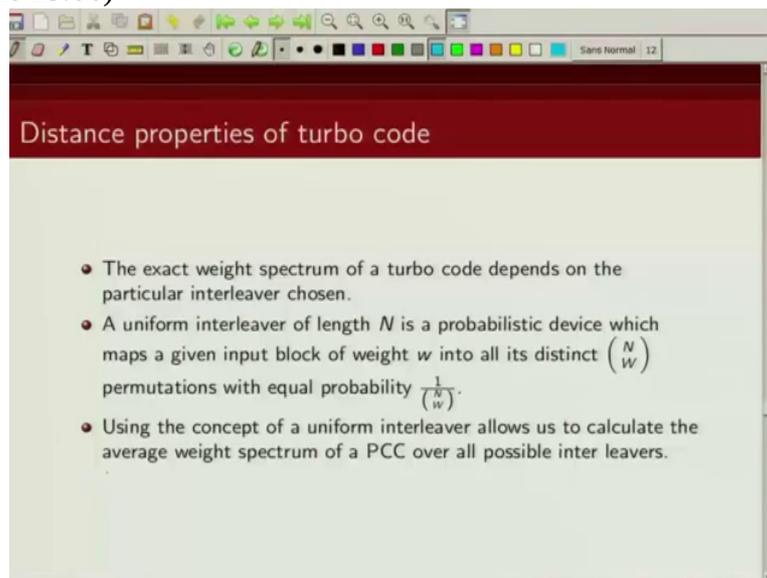
how does uniform interleaver concept helps? Now we can come up with weight distribution

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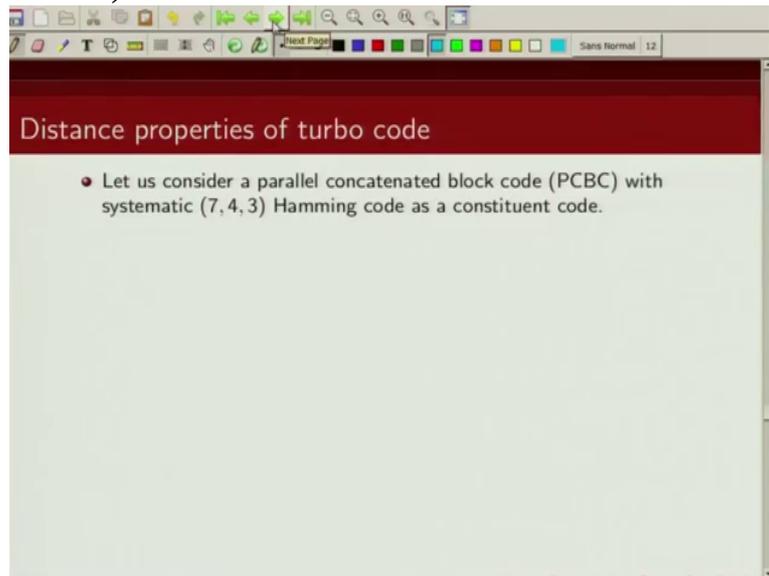
averaged over all possible interleavers and this helps us to compute the

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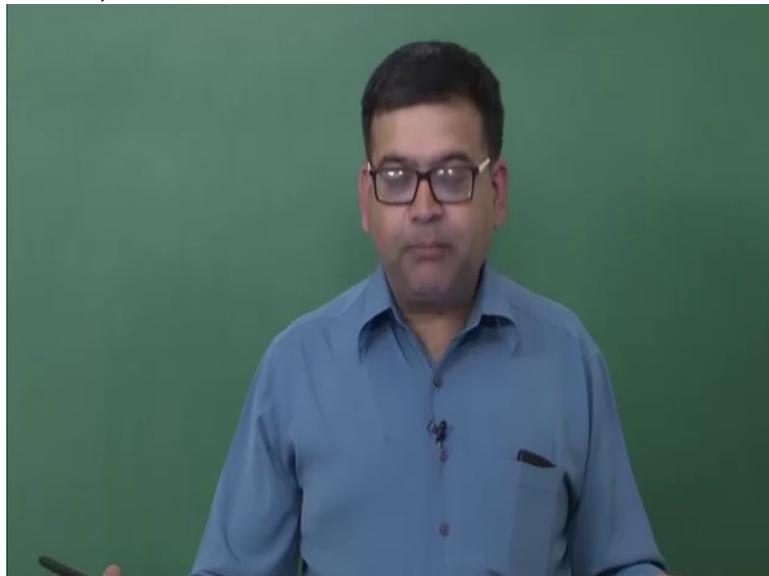
average weight spectrum of parallel concatenated codes.

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So again we are going to take a very simple example. This time we are going to use our favorite 7 4 Hamming codes as constituent encoders so we are going to use a 7 4 Hamming codes as constituent encoder for turbo code and both these constituent encoders use 7 4 Hamming code. So we are basically constructing a

(Refer Slide Time 23:36)



parallel concatenated block code.

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Distance properties of turbo code

- Let us consider a parallel concatenated block code (PCBC) with systematic (7, 4, 3) Hamming code as a constituent code.
- The weight enumerating function (WEF) of this code is given by

$$A(X) = \sum_d A_d X^d = 7X^3 + 7X^4 + X^7$$

Now we know in systematic form the weight enumerating function of Hamming code is given by this expression. There are 7 codewords of weight 3, 7 codewords of weight 4 and 1 codeword of weight 7. These are the

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Distance properties of turbo code

- Let us consider a parallel concatenated block code (PCBC) with systematic (7, 4, 3) Hamming code as a constituent code.
- The weight enumerating function (WEF) of this code is given by

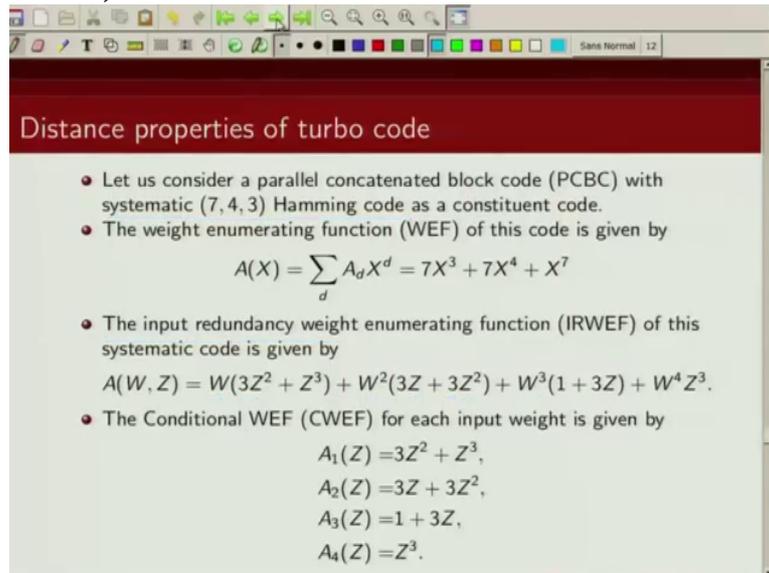
$$A(X) = \sum_d A_d X^d = 7X^3 + 7X^4 + X^7$$

- The input redundancy weight enumerating function (IRWEF) of this systematic code is given by

$$A(W, Z) = W(3Z^2 + Z^3) + W^2(3Z + 3Z^2) + W^3(1 + 3Z) + W^4 Z^3.$$

non zero codewords in a systematic 7 4 Hamming code. Now the input redundancy weight enumerating function, this we have already defined when we talked about weight distribution of convolutional code. So the input redundancy weight enumerating function for this systematic Hamming code is given by this, so for our information sequence of weight 1, there exists 3 sequences which are parity weight 2 and one which have parity weight 3. For information sequence of weight 2, there are three sequences which are parity weight 1, three sequences which are parity weight 2 and so on. Now we can also find out the

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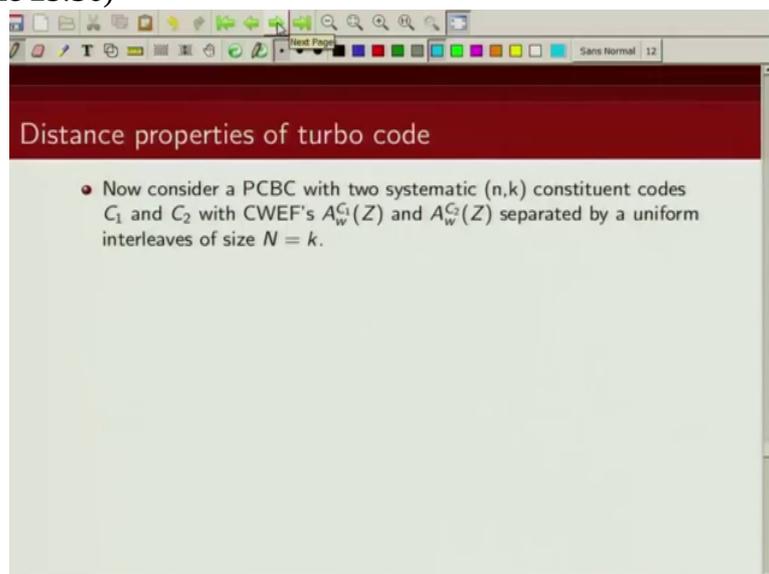


The slide is titled "Distance properties of turbo code" and contains the following content:

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- The input redundancy weight enumerating function (IRWEF) of this systematic code is given by
$$A(W, Z) = W(3Z^2 + Z^3) + W^2(3Z + 3Z^2) + W^3(1 + 3Z) + W^4 Z^3.$$
- The Conditional WEF (CWEF) for each input weight is given by
$$A_1(Z) = 3Z^2 + Z^3,$$
$$A_2(Z) = 3Z + 3Z^2,$$
$$A_3(Z) = 1 + 3Z,$$
$$A_4(Z) = Z^3.$$

conditional weight enumerating function. So from this input redundancy weight enumerating function we can find out conditional weight enumerating function. So when input is 1, this is what we get. So the conditional weight enumerating function is given by, for input 1, is 3 z square plus z cube. Similarly for input weight 2, the conditional weight enumerating function is given by this, so that's this. Similarly for weight 3, the conditionally weight enumerating function is 1 plus 3 z and for input weight 4, the conditional weight enumerating function is given by z cube.

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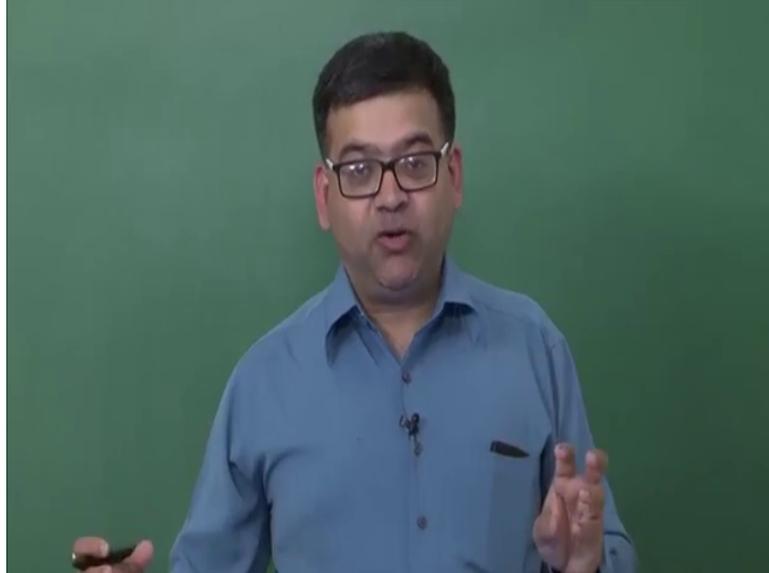
The slide is titled "Distance properties of turbo code" and contains the following content:

- Now consider a PCBC with two systematic (n,k) constituent codes C_1 and C_2 with CWEF's $A_{W_1}^{C_1}(Z)$ and $A_{W_2}^{C_2}(Z)$ separated by a uniform interleaves of size $N = k$.

Now let us consider a parallel concatenation of these two Hamming codes, Ok. So we have these 2 systematic Hamming codes c_1, c_2 whose conditional weight enumerating function is

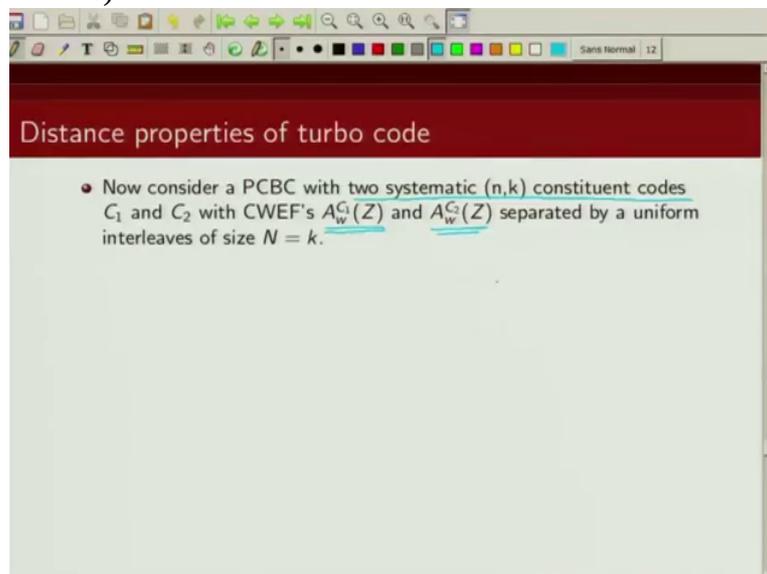
given by this and this. And we apply this concept of uniform interleaver to find out what is the conditional weight

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enumerating function of the parallel concatenated code. Now again recall

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our encoder is something like this. So we have systematic encoder and then we have information bit which is interleaved and fed to second encoder. This is the output.

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Distance properties of turbo code

- Now consider a PCBC with two systematic (n,k) constituent codes C_1 and C_2 with CWF's $A_w^{C_1}(Z)$ and $A_w^{C_2}(Z)$ separated by a uniform interleaves of size $N = k$.

The diagram illustrates a parallel concatenated code (PCBC) structure. It shows two constituent codes, C_1 and C_2 , each represented by a rectangular block. The input to the first encoder is a sequence of bits, and the output is a sequence of bits. The output of the first encoder is interleaved with the input of the second encoder. The interleaving is done by a uniform interleaver of size $N = k$. The output of the second encoder is a sequence of bits. The overall structure is a parallel concatenated code.

So

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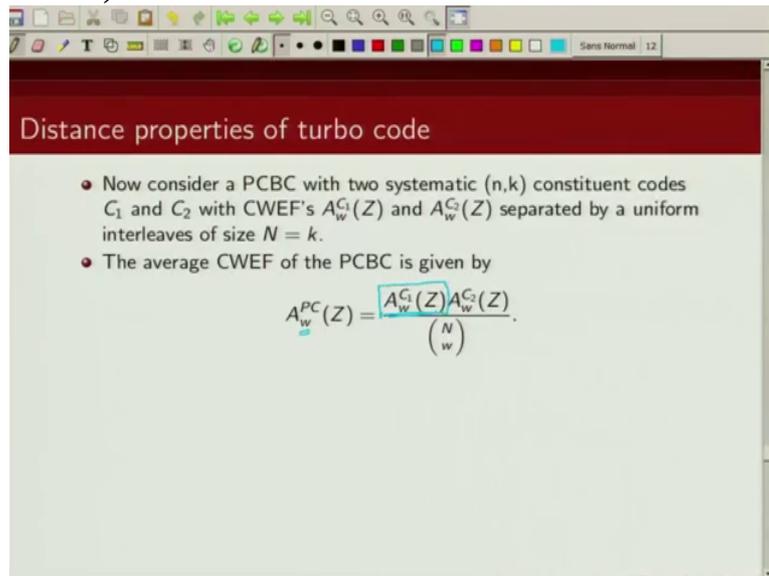
Distance properties of turbo code

- Now consider a PCBC with two systematic (n,k) constituent codes C_1 and C_2 with CWF's $A_w^{C_1}(Z)$ and $A_w^{C_2}(Z)$ separated by a uniform interleaves of size $N = k$.
- The average CWF of the PCBC is given by

$$A_w^{PC}(Z) = \frac{A_w^{C_1}(Z)A_w^{C_2}(Z)}{\binom{N}{w}}$$

conditional weight enumerating function of the parallel concatenated code can be given by this, so for first encoder we have input w getting in and this will give me the conditional

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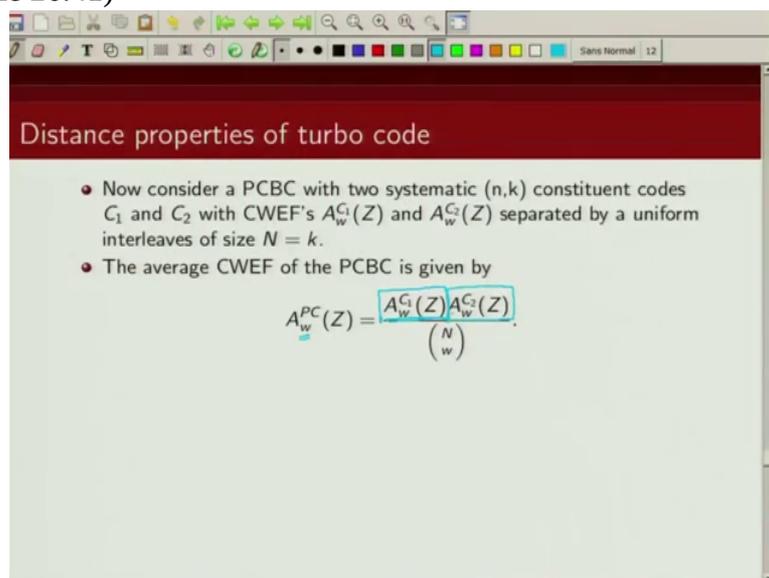
The slide is titled "Distance properties of turbo code" and contains the following text:

- Now consider a PCBC with two systematic (n,k) constituent codes C_1 and C_2 with CWEF's $A_w^{C_1}(Z)$ and $A_w^{C_2}(Z)$ separated by a uniform interleaves of size $N = k$.
- The average CWEF of the PCBC is given by

$$A_w^{PC}(Z) = \frac{A_w^{C_1}(Z)A_w^{C_2}(Z)}{\binom{N}{w}}$$

weight enumerating function corresponding to input weight w . Similarly for the second encoder

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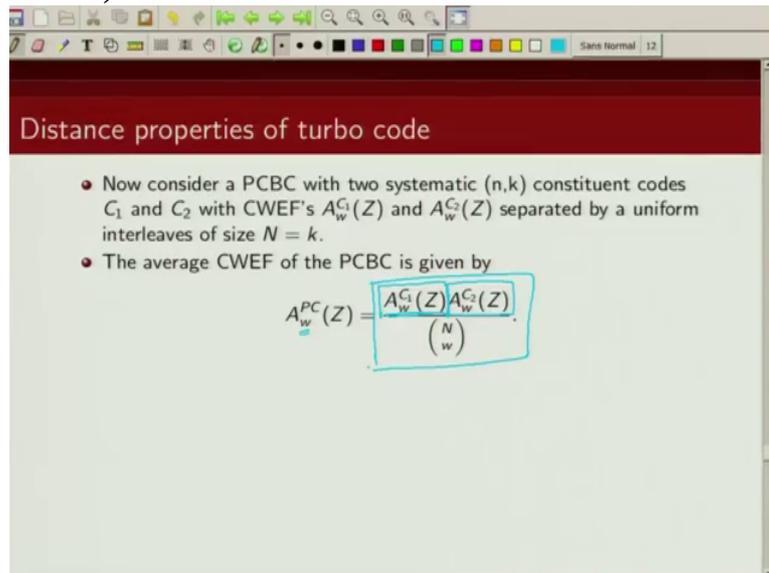
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- The average CWEF of the PCBC is given by

$$A_w^{PC}(Z) = \frac{A_w^{C_1}(Z)A_w^{C_2}(Z)}{\binom{N}{w}}$$

this will be the weight enumerating function for input w and then averaging it out over all

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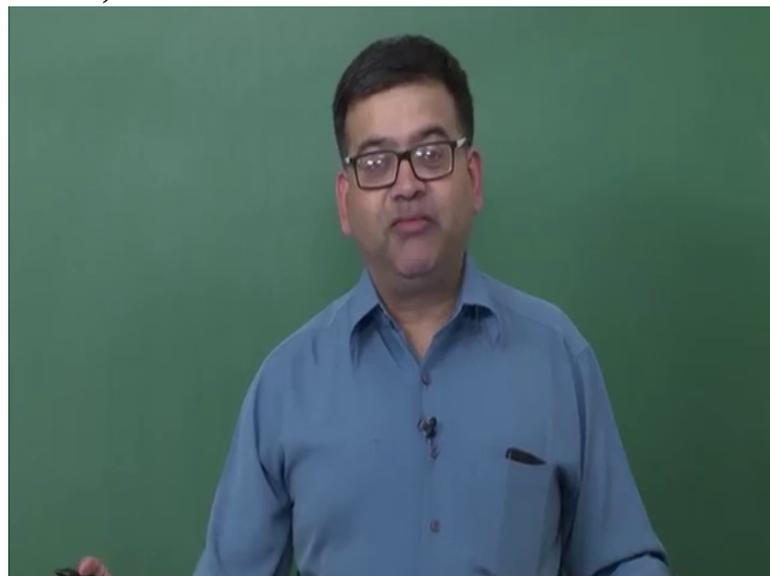
Distance properties of turbo code

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- The average CWEF of the PCBC is given by

$$A_w^{PC}(Z) = \frac{A_w^{C_1}(Z) A_w^{C_2}(Z)}{\binom{N}{w}}$$

possible interleavers I get this average conditional weight enumerating function given by this expression. So the nice thing about this uniform

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interleaver concept is it allows us to compute the average weight enumerating function of this parallel concatenated code which is independent of the type of interleaver that we are using.

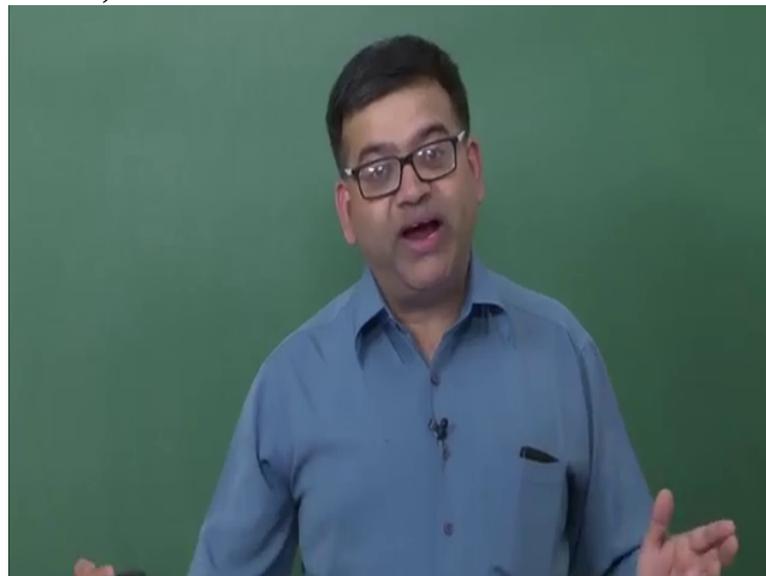
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Distance properties of turbo code

- Now consider a PCBC with two systematic (n,k) constituent codes C_1 and C_2 with CWEF's $A_w^{C_1}(Z)$ and $A_w^{C_2}(Z)$ separated by a uniform interleaves of size $N = k$.
- The average CWEF of the PCBC is given by
$$A_w^{PC}(Z) = \frac{A_w^{C_1}(Z)A_w^{C_2}(Z)}{\binom{N}{w}}$$
- The average IRWEF is given by
$$A_w^{PC}(W, Z) = \sum_{(1 \leq w \leq N)} W^w A_w^{PC}(Z)$$

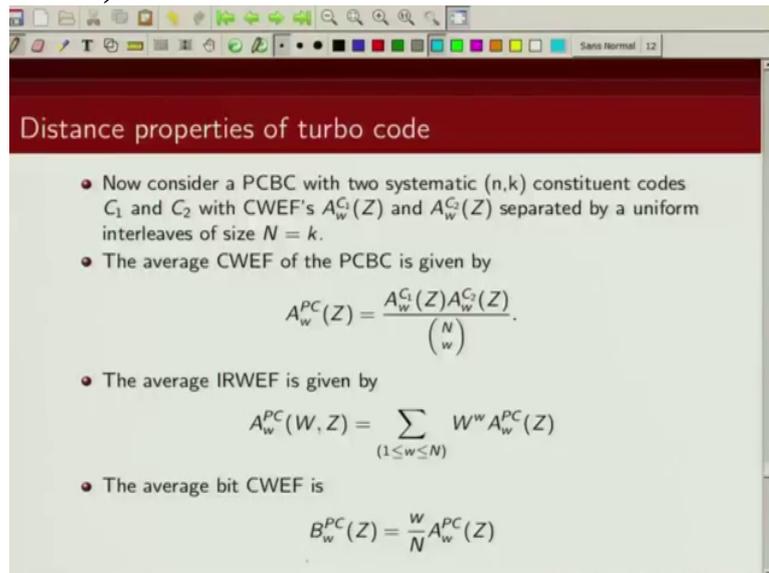
Now similarly we could define average input redundancy weight enumerating function is given by this expression. Now please note this is a, these are straightforward extensions of the definitions that we have studied when talked about

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weight enumerating function of convolutional code.

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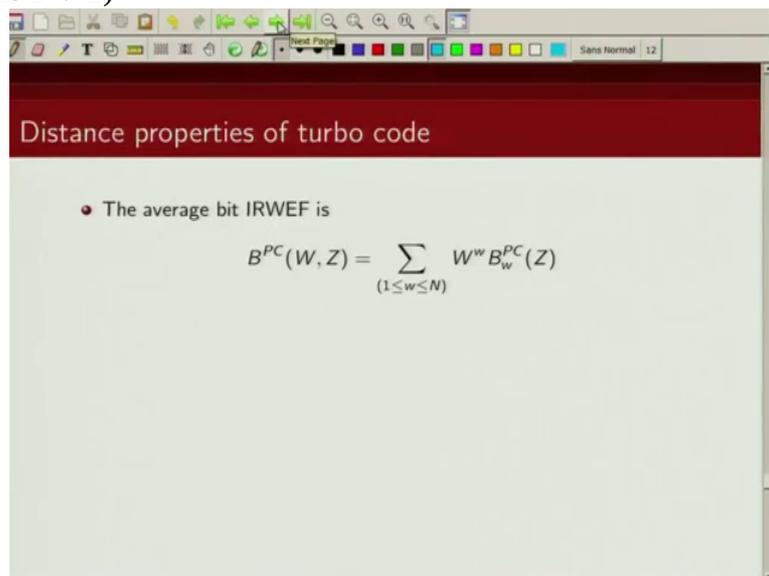


The slide is titled "Distance properties of turbo code" and contains the following text and equations:

- Now consider a PCBC with two systematic (n,k) constituent codes C_1 and C_2 with CWEF's $A_w^{C_1}(Z)$ and $A_w^{C_2}(Z)$ separated by a uniform interleaves of size $N = k$.
- The average CWEF of the PCBC is given by
$$A_w^{PC}(Z) = \frac{A_w^{C_1}(Z)A_w^{C_2}(Z)}{\binom{N}{w}}$$
- The average IRWEF is given by
$$A_w^{PC}(W, Z) = \sum_{(1 \leq w \leq N)} W^w A_w^{PC}(Z)$$
- The average bit CWEF is
$$B_w^{PC}(Z) = \frac{w}{N} A_w^{PC}(Z)$$

We can define bit conditional weight enumerating function,

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The slide is titled "Distance properties of turbo code" and contains the following text and equation:

- The average bit IRWEF is
$$B^{PC}(W, Z) = \sum_{(1 \leq w \leq N)} W^w B_w^{PC}(Z)$$

average bit input redundancy weight enumerating function,

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Distance properties of turbo code

- The average bit IRWEF is

$$B^{PC}(W, Z) = \sum_{(1 \leq w \leq N)} W^w B_w^{PC}(Z)$$

- The average codeword WEF is

$$A_w^{PC}(X) = \sum_{d_{min} \leq d} A_d X^d = \sum_{1 \leq w \leq N} W^w A_w^{PC}(Z)|_{W=Z=X}$$

average codeword weight enumerating function and

(Refer Slide Time 27:55)

Distance properties of turbo code

- The average bit IRWEF is

$$B^{PC}(W, Z) = \sum_{(1 \leq w \leq N)} W^w B_w^{PC}(Z)$$

- The average codeword WEF is

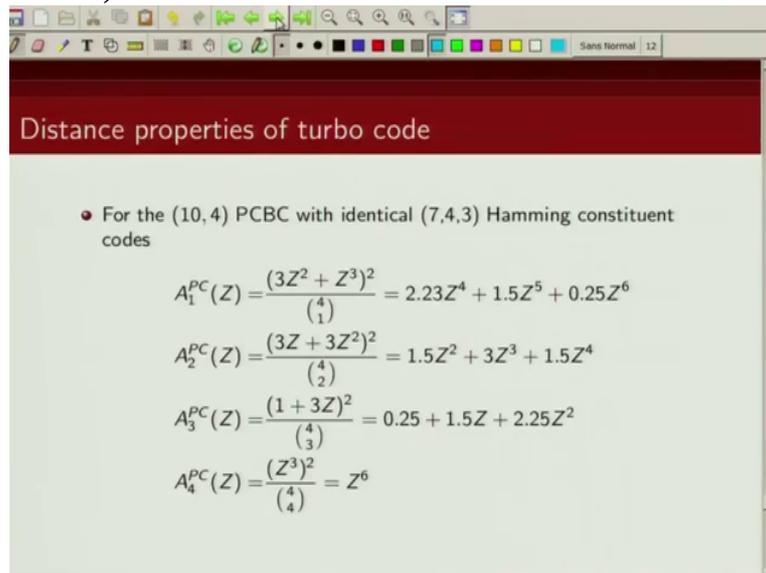
$$A_w^{PC}(X) = \sum_{d_{min} \leq d} A_d X^d = \sum_{1 \leq w \leq N} W^w A_w^{PC}(Z)|_{W=Z=X}$$

- The average bit WEF is

$$B_w^{PC}(X) = \sum_{d_{min} \leq d} B_d X^d = \sum_{1 \leq w \leq N} W^w B_w^{PC}(Z)|_{W=Z=X}$$

average bit weight enumerating function .

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Distance properties of turbo code

- For the (10, 4) PCBC with identical (7,4,3) Hamming constituent codes

$$A_1^{PC}(Z) = \frac{(3Z^2 + Z^3)^2}{\binom{4}{1}} = 2.23Z^4 + 1.5Z^5 + 0.25Z^6$$

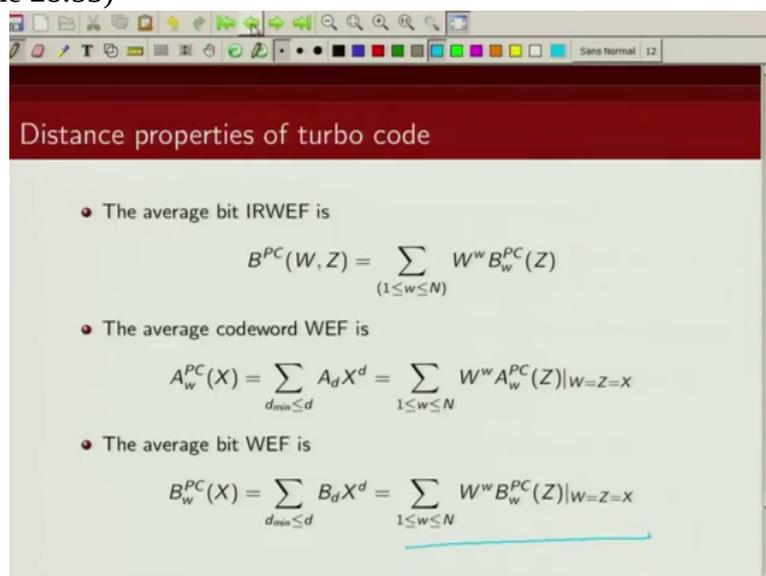
$$A_2^{PC}(Z) = \frac{(3Z + 3Z^2)^2}{\binom{4}{2}} = 1.5Z^2 + 3Z^3 + 1.5Z^4$$

$$A_3^{PC}(Z) = \frac{(1 + 3Z)^2}{\binom{4}{3}} = 0.25 + 1.5Z + 2.25Z^2$$

$$A_4^{PC}(Z) = \frac{(Z^3)^2}{\binom{4}{4}} = Z^6$$

Now let us take an example to illustrate this. So we have this parallel concatenated code using 7 4 Hamming code. So input is length 4 and we get 3 parity bits from first encoder, we get 3 parity bits from the second encoder and of course there are 4 systematic bits, so overall codeword length is 10. So what we get is a 10 4 parallel concatenated block code which I am using identically considering encoders and there is the odd 7 4 Hamming code. Now using the expressions I have said before, we already know the conditional weight enumerating function of the constituent encoder. Now using this

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Distance properties of turbo code

- The average bit IRWEF is

$$B^{PC}(W, Z) = \sum_{(1 \leq w \leq N)} W^w B_w^{PC}(Z)$$

- The average codeword WEF is

$$A_w^{PC}(X) = \sum_{d_{min} \leq d} A_d X^d = \sum_{1 \leq w \leq N} W^w A_w^{PC}(Z)|_{W=Z=X}$$

- The average bit WEF is

$$B_w^{PC}(X) = \sum_{d_{min} \leq d} B_d X^d = \sum_{1 \leq w \leq N} W^w B_w^{PC}(Z)|_{W=Z=X}$$

expression,

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The slide is titled "Distance properties of turbo code" and contains the following text and equations:

- Now consider a PCBC with two systematic (n,k) constituent codes C_1 and C_2 with CWEF's $A_w^{C_1}(Z)$ and $A_w^{C_2}(Z)$ separated by a uniform interleaves of size $N = k$.
- The average CWEF of the PCBC is given by
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- The average bit CWEF is
$$B_w^{PC}(Z) = \frac{w}{N} A_w^{PC}(Z)$$

I can compute

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The slide is titled "Distance properties of turbo code" and contains the following text and equations:

- Now consider a PCBC with two systematic (n,k) constituent codes C_1 and C_2 with CWEF's $A_w^{C_1}(Z)$ and $A_w^{C_2}(Z)$ separated by a uniform interleaves of size $N = k$.
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- The average bit CWEF is
$$B_w^{PC}(Z) = \frac{w}{N} A_w^{PC}(Z)$$

the conditional weight enumerating function of the parallel concatenated block code. So these are the

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Distance properties of turbo code

- For the (10,4) PCBC with identical (7,4,3) Hamming constituent codes

$$A_1^{PC}(Z) = \frac{(3Z^2 + Z^3)^2}{\binom{4}{1}} = 2.23Z^4 + 1.5Z^5 + 0.25Z^6$$

$$A_2^{PC}(Z) = \frac{(3Z + 3Z^2)^2}{\binom{4}{2}} = 1.5Z^2 + 3Z^3 + 1.5Z^4$$

$$A_3^{PC}(Z) = \frac{(1 + 3Z)^2}{\binom{4}{3}} = 0.25 + 1.5Z + 2.25Z^2$$

$$A_4^{PC}(Z) = \frac{(Z^3)^2}{\binom{4}{4}} = Z^6$$

average weight enumerating, so corresponding to input 1, so you can there are 2 point 2 3 on average sequences of parity weight 4, 1 point 5 sequences of parity weight 5 and point 2 5 on an average sequences of parity weight 6. Similarly we could find out for input weight 2, input weight 3, input weight 4. Now note the difference from the constituent encoders that we have studied. Here you can see these numbers, its coefficients are coming as non integers,

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Distance properties of turbo code

- For the (10,4) PCBC with identical (7,4,3) Hamming constituent codes

$$A_1^{PC}(Z) = \frac{(3Z^2 + Z^3)^2}{\binom{4}{1}} = 2.23Z^4 + 1.5Z^5 + 0.25Z^6$$

$$A_2^{PC}(Z) = \frac{(3Z + 3Z^2)^2}{\binom{4}{2}} = 1.5Z^2 + 3Z^3 + 1.5Z^4$$

$$A_3^{PC}(Z) = \frac{(1 + 3Z)^2}{\binom{4}{3}} = 0.25 + 1.5Z + 2.25Z^2$$

$$A_4^{PC}(Z) = \frac{(Z^3)^2}{\binom{4}{4}} = Z^6$$

can come as non integers, right? That's because we are averaging over all possible interleavers.

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Distance properties of turbo code

- Similarly, the IRWEF's are given by

$$A^{PC}(W, Z) = W(2.25Z^4 + 1.5Z^5 + 0.25Z^6) + W^2(1.5Z^2 + 3Z^3 + 1.5Z^4) + W^3(0.25 + 1.5Z + 2.25Z^2) + W^4Z^6$$

$$B^{PC}(W, Z) = W(0.56Z^4 + 0.37Z^5 + 0.06Z^6) + W^2(0.75Z^2 + 1.5Z^3 + 0.75Z^4) + W^3(0.19 + 1.12 + 1.69Z^2) + W^4Z^6$$

And again following the definition of input redundancy weight enumerating function we can find out the input redundancy weight enumerating function, bit input redundancy weight enumerating function which is given by these two expressions. And

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Distance properties of turbo code

- Similarly, the IRWEF's are given by

$$A^{PC}(W, Z) = W(2.25Z^4 + 1.5Z^5 + 0.25Z^6) + W^2(1.5Z^2 + 3Z^3 + 1.5Z^4) + W^3(0.25 + 1.5Z + 2.25Z^2) + W^4Z^6$$

$$B^{PC}(W, Z) = W(0.56Z^4 + 0.37Z^5 + 0.06Z^6) + W^2(0.75Z^2 + 1.5Z^3 + 0.75Z^4) + W^3(0.19 + 1.12 + 1.69Z^2) + W^4Z^6$$

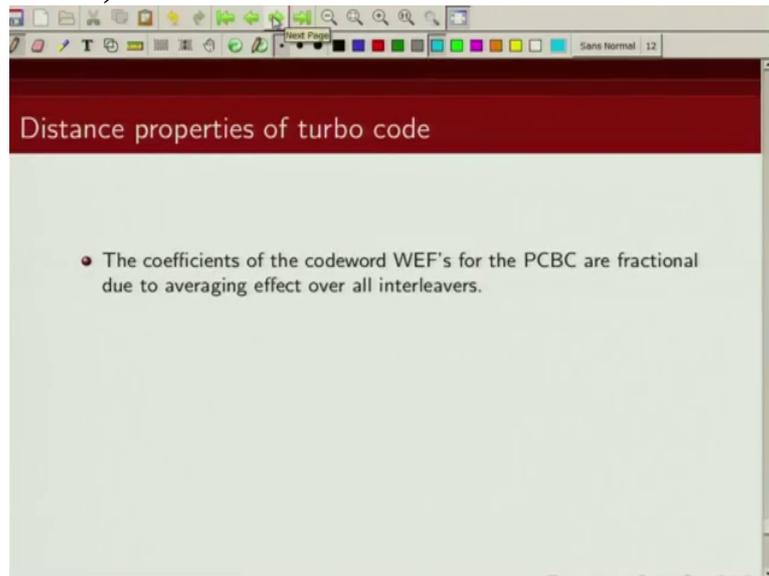
- Also, WEF's are given by

$$A^{PC}(X) = 0.25X^3 + 3X^4 + 7.5X^5 + 3X^6 + 0.25X^7 + X^{10}$$

$$B^{PC}(X) = 0.19X^3 + 1.87X^4 + 3.75X^5 + 1.12X^6 + 0.06X^7 + X^{10}$$

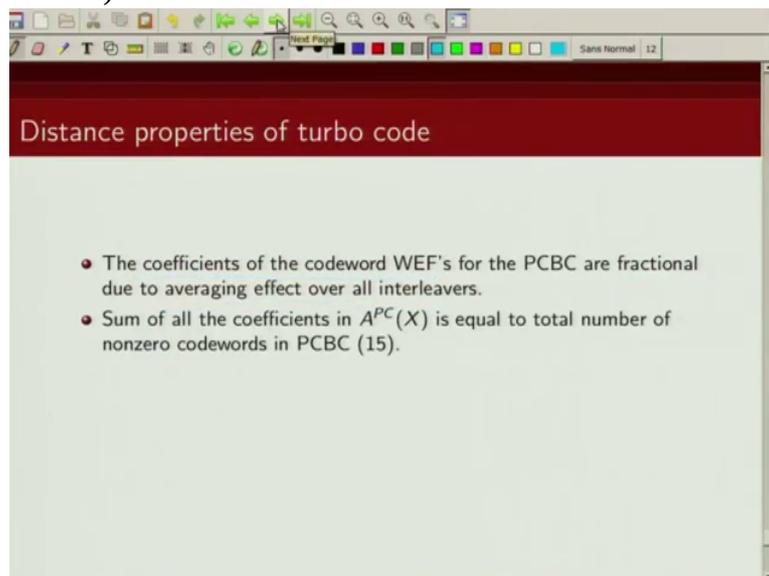
similarly if we put w equal to, z is equal to x, we can get the weight enumerating function and the bit weight enumerating function which is given by this. Now points to be noted.

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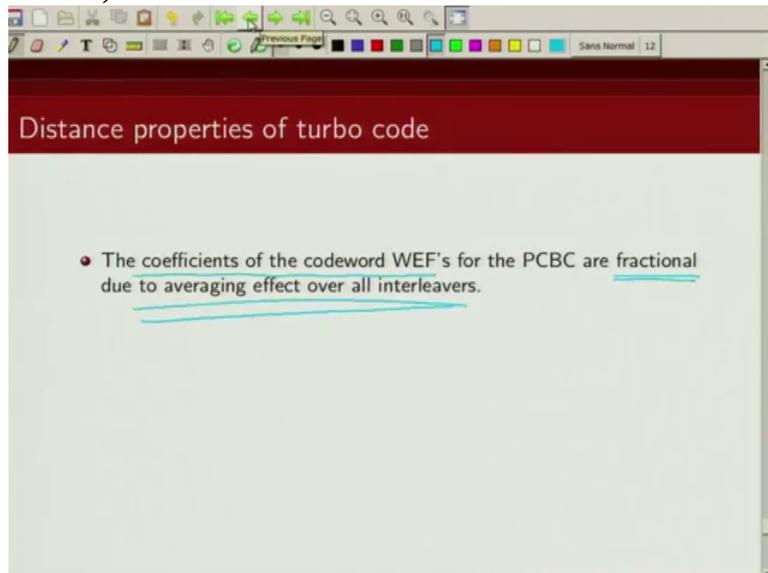
We saw the coefficients of the codeword weight enumerating function are fractional and why they are fractional, because this is due to averaging effect over all interleavers. However

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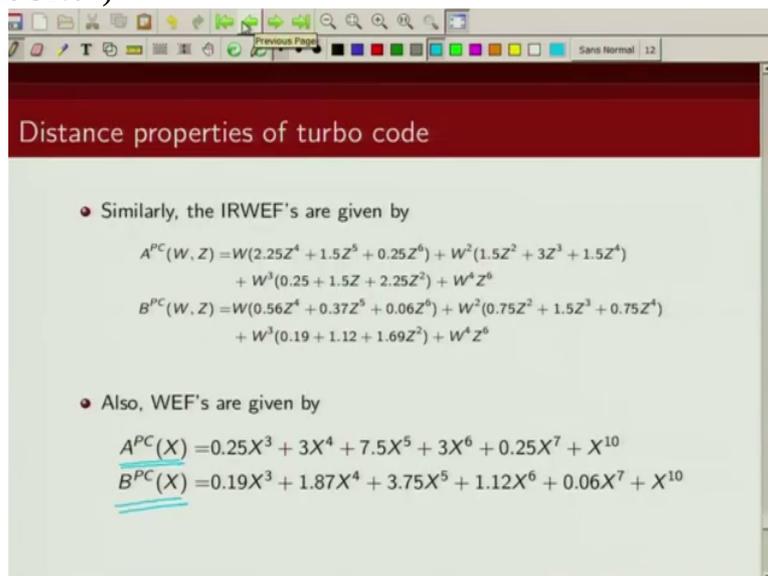
if you sum up all the coefficients, they add up to total number of non zero codewords. You can

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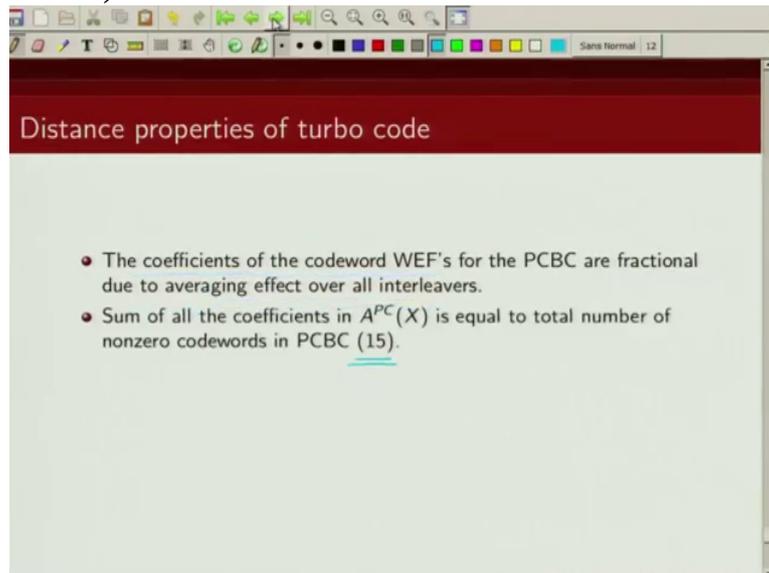
verify that. If you

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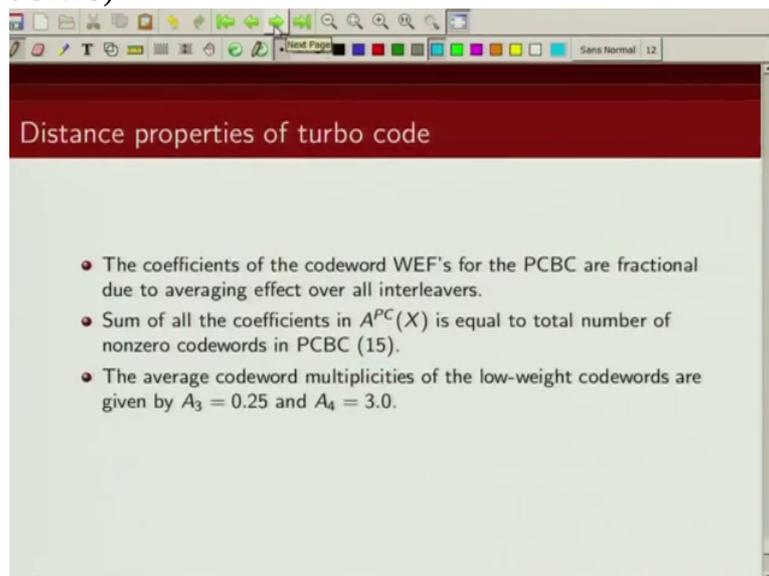
add up this, Ok, point 2 5, point 2 5 is point 5, 7 point 5 this is 8, 8 plus 6 is 14 plus 1 15. So you can see, if you add

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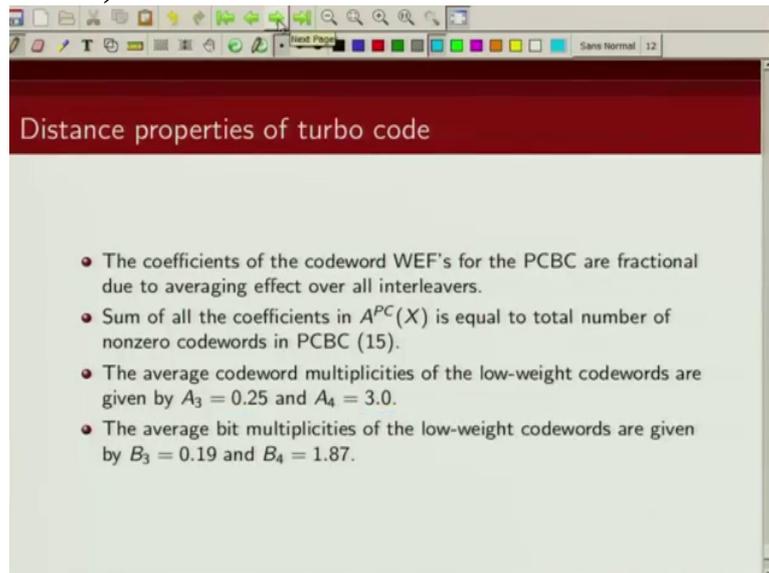
them up we get the total number of non zero codewords.

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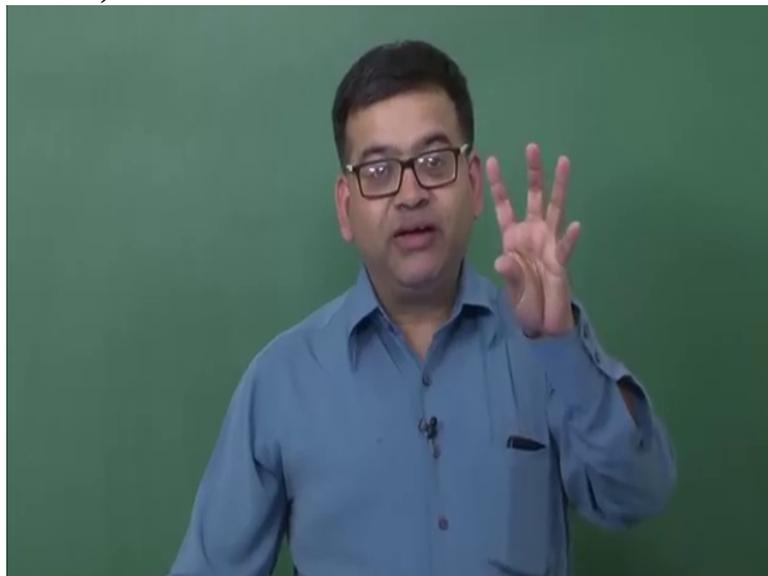
And you can also see that average codeword multiplicity has reduced. Its point 2 5 when the input, corresponding to input of 3 and it is 3 corresponding to input of 4,

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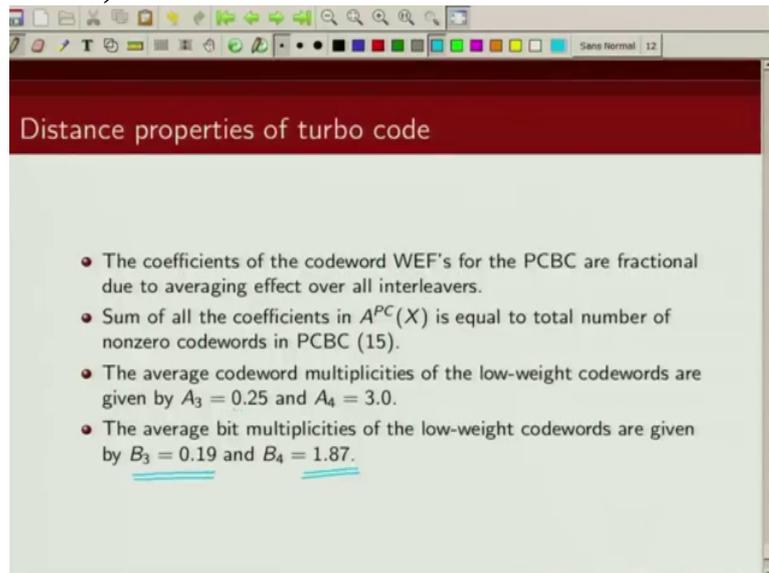
similarly bit multiplicity for input of 3 is point 1 9 and it is 1 point 8 7 for input of length, of weight 4. So if you compare this code multiplicity from the original code that we used, we were using

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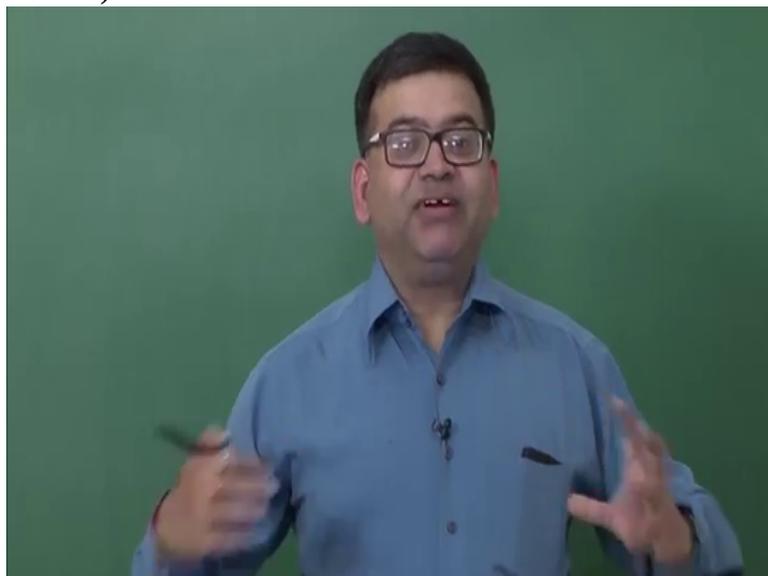
7 4 Hamming code which had 7 codewords of weight 3. Here you can see

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the weight multiplicity is point 2.5. So you can see the effect

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of interleaver on code multiplicity and you can use this concept of uniform interleaver to compute your average weight enumerating function, average input output weight enumerating function, average bit conditional weight enumerating function. So all those quantities that we have defined in the context of weight enumerating function for convolutional code, you could do similar analysis for these parallel concatenated codes using this idea of uniform interleaver. So with this I am going to conclude the discussion on these distance properties of turbo codes, thank you.